

Nonlinear Optimization

**Part III**  
**Numerical algorithms**

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# Unconstrained minimization

Optimization problem:

$$\text{minimize } f(x)$$

Assumptions:

- $f : \mathbb{R}^n \rightarrow \mathbb{R} \cup \{+\infty\}$  is convex
- $\text{dom } f = \{x : f(x) < +\infty\}$  is open
- $f$  is closed
- $f$  is twice continuously differentiable
- problem is solvable:  $\exists x^* : f(x^*) = p^* := \inf\{f(x) : x\}$

**Example:** quadratic ( $A \succeq 0$ )

$$\text{minimize } x^\top Ax + b^\top x + c$$

- $\text{dom } f = \mathbb{R}^n$
- closed
- bounded below iff  $b \in \text{span}(A)$

**Example:** log-sum-exp

$$\text{minimize} \quad \log \left( \sum_{i=1}^m e^{a_i^\top x + b_i} \right)$$

- $\text{dom } f = \mathbb{R}^n$
- closed
- bounded below iff  $0 \in \text{co} \{a_i\}$

**Example:** log barrier for polyhedron

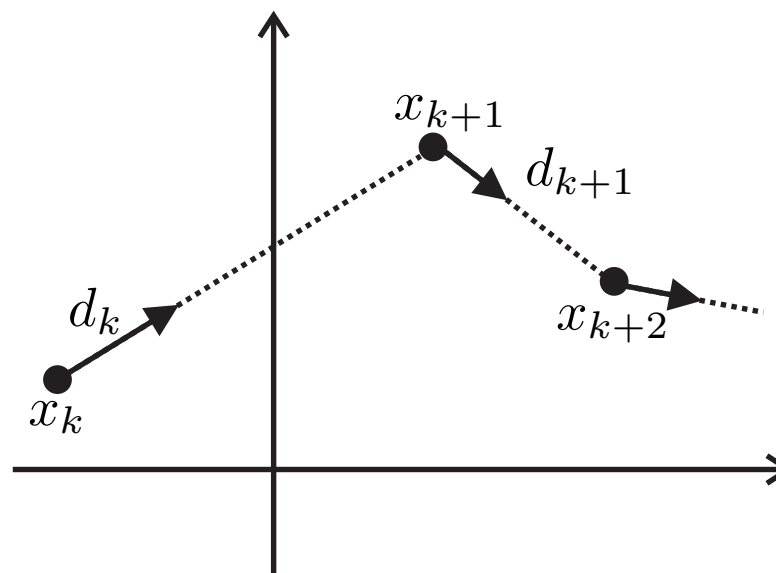
$$\text{minimize} \quad \sum_{i=1}^m -\log(b_i - a_i^\top x)$$

- $\text{dom } f = \{x : Ax < b\}$
- closed
- bounded below if there exists  $\lambda > 0$ ,  $\sum_k \lambda_k a_k = 0$

Algorithm produces a minimizing sequence:  $x_1, x_2, x_3, \dots$  such that

$$f(x_k) \downarrow p^\star$$

Basic iteration of line-search algorithms:  $x_{k+1} = x_k + t_k d_k$



- $d_k$  is the search direction
- $t_k > 0$  is the step size

**Definition (Descent direction)**  $d_k$  is a descent direction for  $f$  at a non-stationary point  $x_k \in \text{dom } f$  if

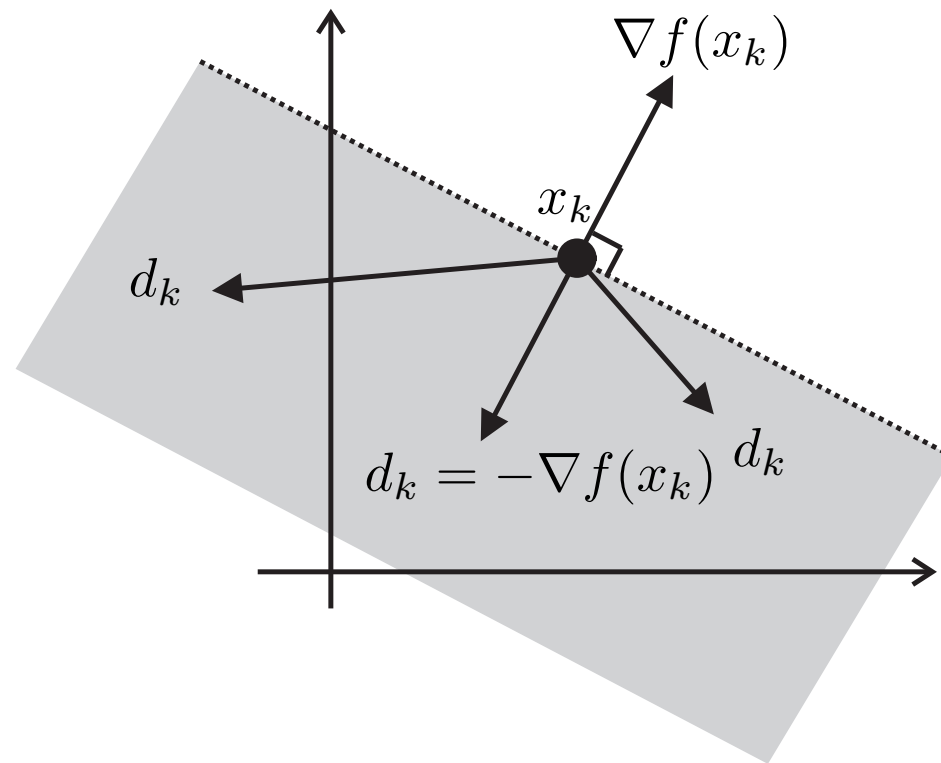
$$\langle \nabla f(x_k), d_k \rangle < 0$$

Example:  $d_k = -\nabla f(x_k)$  is a descent direction because

$$\langle \nabla f(x_k), d_k \rangle = -\|\nabla f(x_k)\|^2 < 0$$



Infinite number of choices for descent directions  $d_k$ :



Important: if  $d_k$  is a descent direction, then

$$\exists \bar{t} > 0 : f(x_k + td_k) < f(x_k) \text{ for all } t \in ]0, \bar{t}]$$

Prototype of a line-search algorithm:

1.     **initialization**   ▶ choose  $x_0 \in \text{dom } f$  and tolerance  $\epsilon > 0$
2.                       ▶ set  $k = 0$
3.     **loop**           ▶ compute  $g_k = \nabla f(x_k)$
4.                       ▶ if  $\|g_k\| < \epsilon$  stop
5.                       ▶ compute descent direction  $d_k$
6.                       ▶ compute step size  $t_k$
7.                       ▶ update  $x_{k+1} = x_k + t_k d_k$
8.                       ▶  $k \leftarrow k + 1$  and return to loop

Line 6 is called the line search subproblem

- exact line search:

$$t_k \in \arg \min_{t \geq 0} \phi(t) := f(x_k + td_k)$$

- Armijo's rule or backtracking line search ( $0 < \alpha < 0.5, 0 < \beta < 1$ ):

$t = 1$  and while  $f(x_k + td_k) > f(x_k) + t\alpha \nabla f(x_k)^\top d_k$  do  $t := \beta t$

**Theorem (Convergence of steepest descent method)** Let  $x_0 \in \text{dom } f$  and suppose there exists  $m > 0$  such that

$$\nabla^2 f(x) \succeq mI \quad \text{for all } x \in S_{f(x_0)} f.$$

Then,

$$x_k \rightarrow x^*.$$

Example:

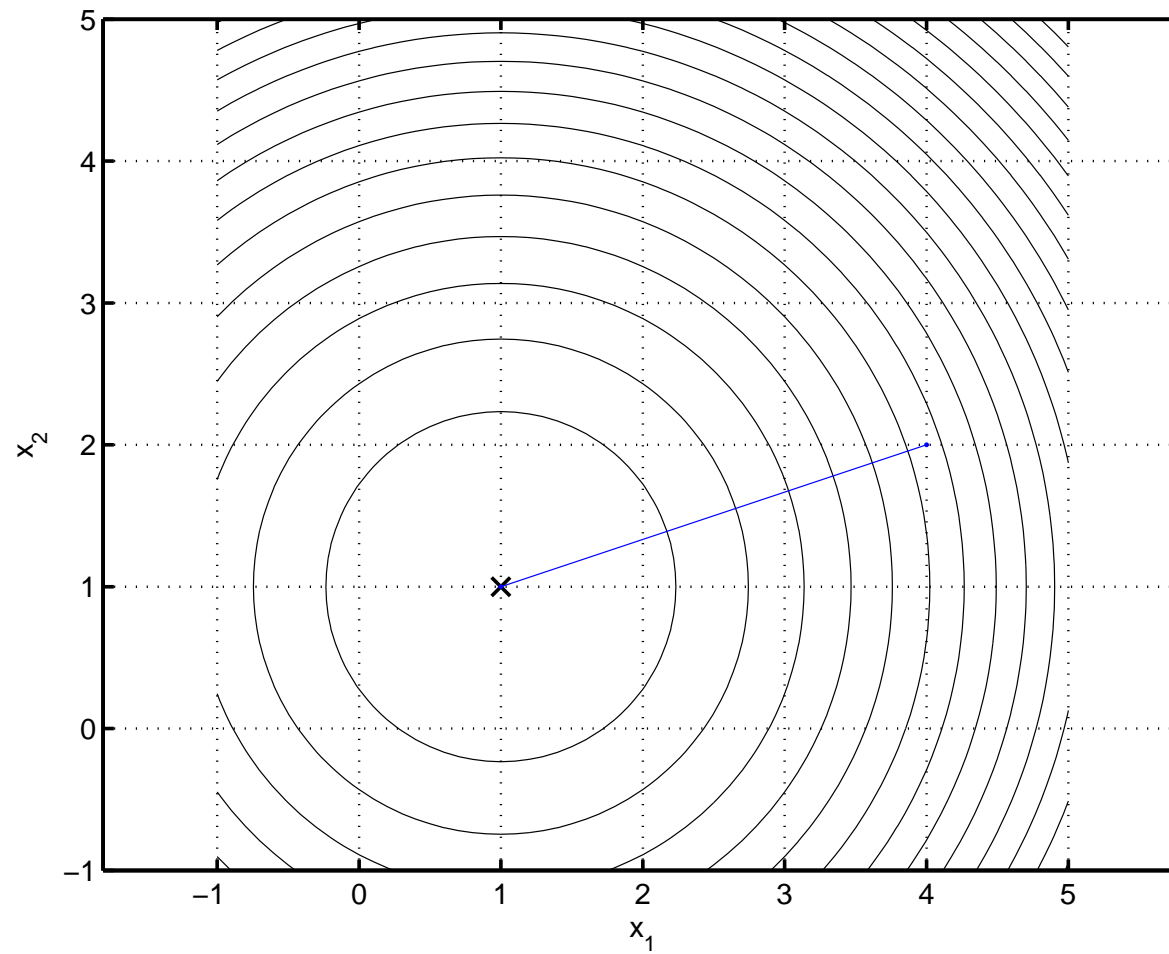
- $f : \mathbb{R}^2 \rightarrow \mathbb{R}$

$$f(x) = \frac{1}{2}(x - c)^\top A(x - c)$$

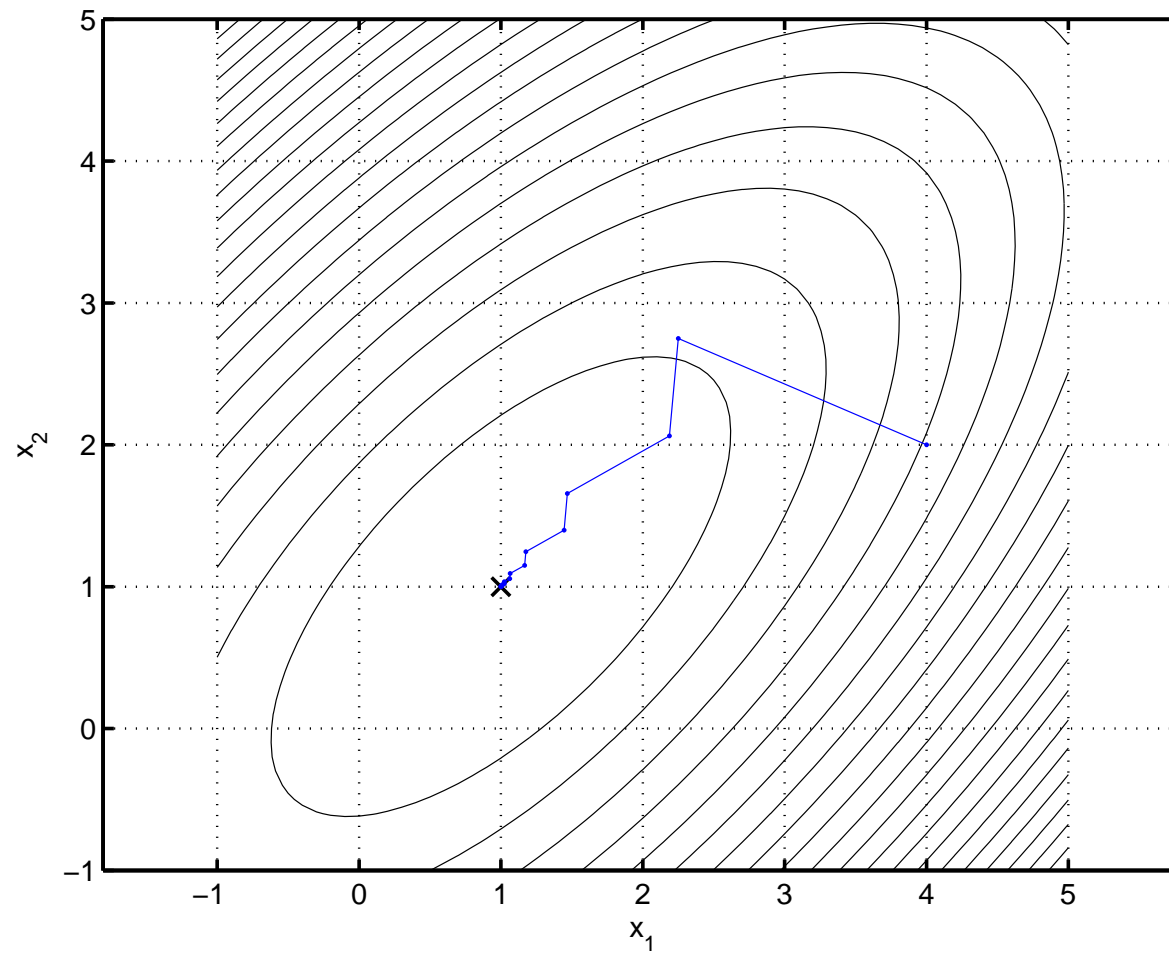
where

$$c = \begin{bmatrix} 1 \\ 1 \end{bmatrix} \quad A = Q \begin{bmatrix} 1 & 0 \\ 0 & \kappa \end{bmatrix} Q^\top \quad Q = \begin{bmatrix} \cos\left(\frac{\pi}{4}\right) & -\sin\left(\frac{\pi}{4}\right) \\ \sin\left(\frac{\pi}{4}\right) & \cos\left(\frac{\pi}{4}\right) \end{bmatrix}$$

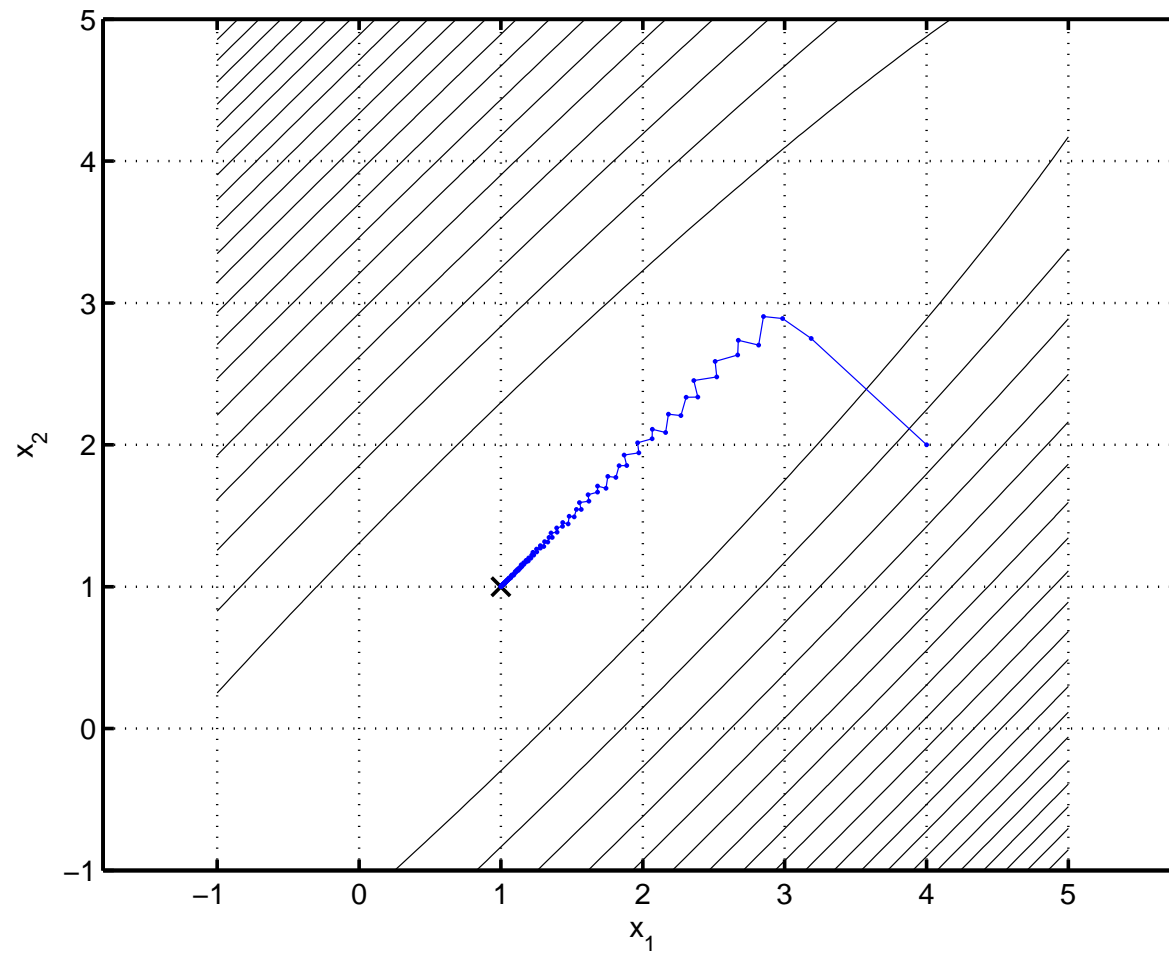
- $x^\star = (1, 1)$  and  $\kappa(\nabla^2 f(x^\star)) = \kappa \geq 1$
- $x_0 = (4, 2)$ ,  $\epsilon = 10^{-6}$ ,  $\alpha = 10^{-4}$ ,  $\beta = 0.5$



$\kappa = 1$  (number of iterations: 1)

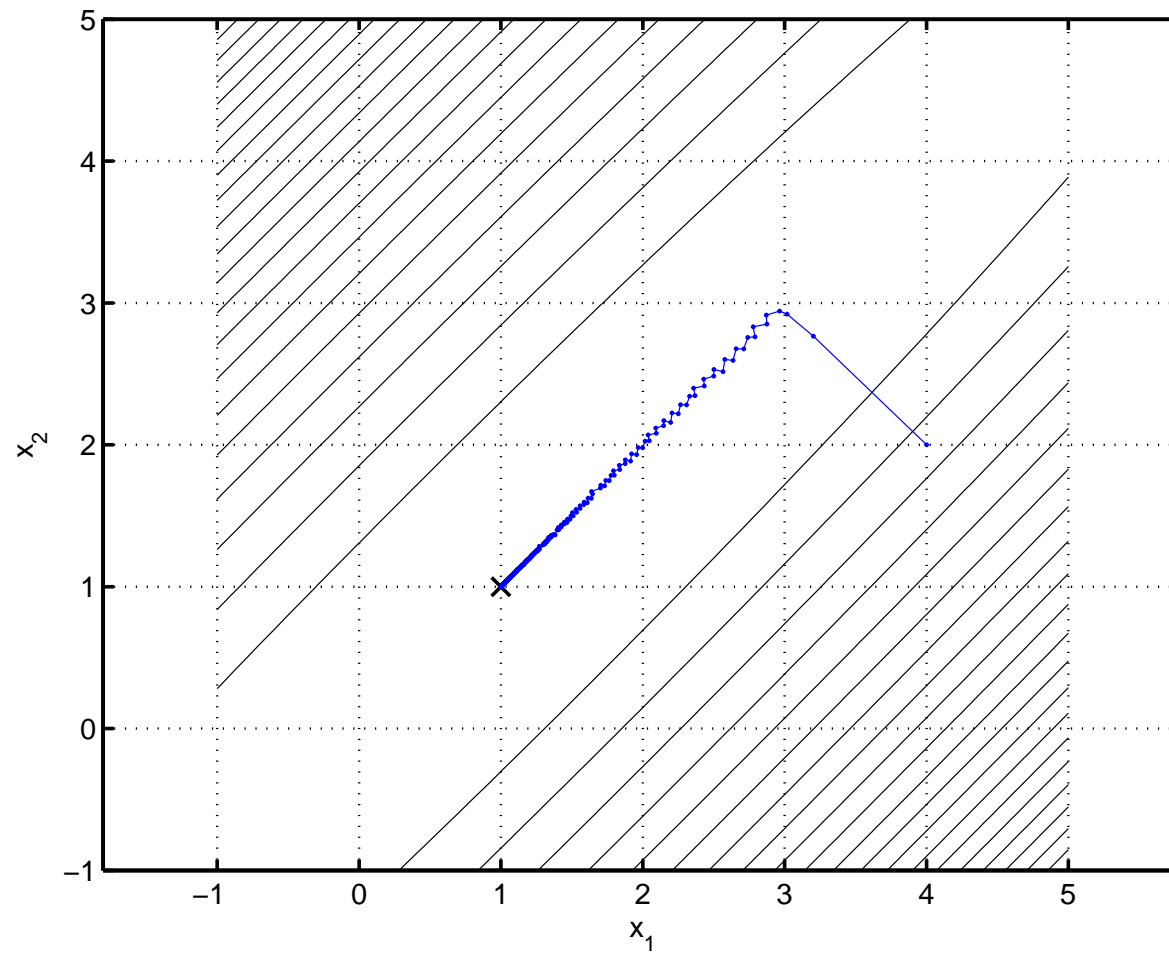


$\kappa = 5$  (number of iterations: 32)



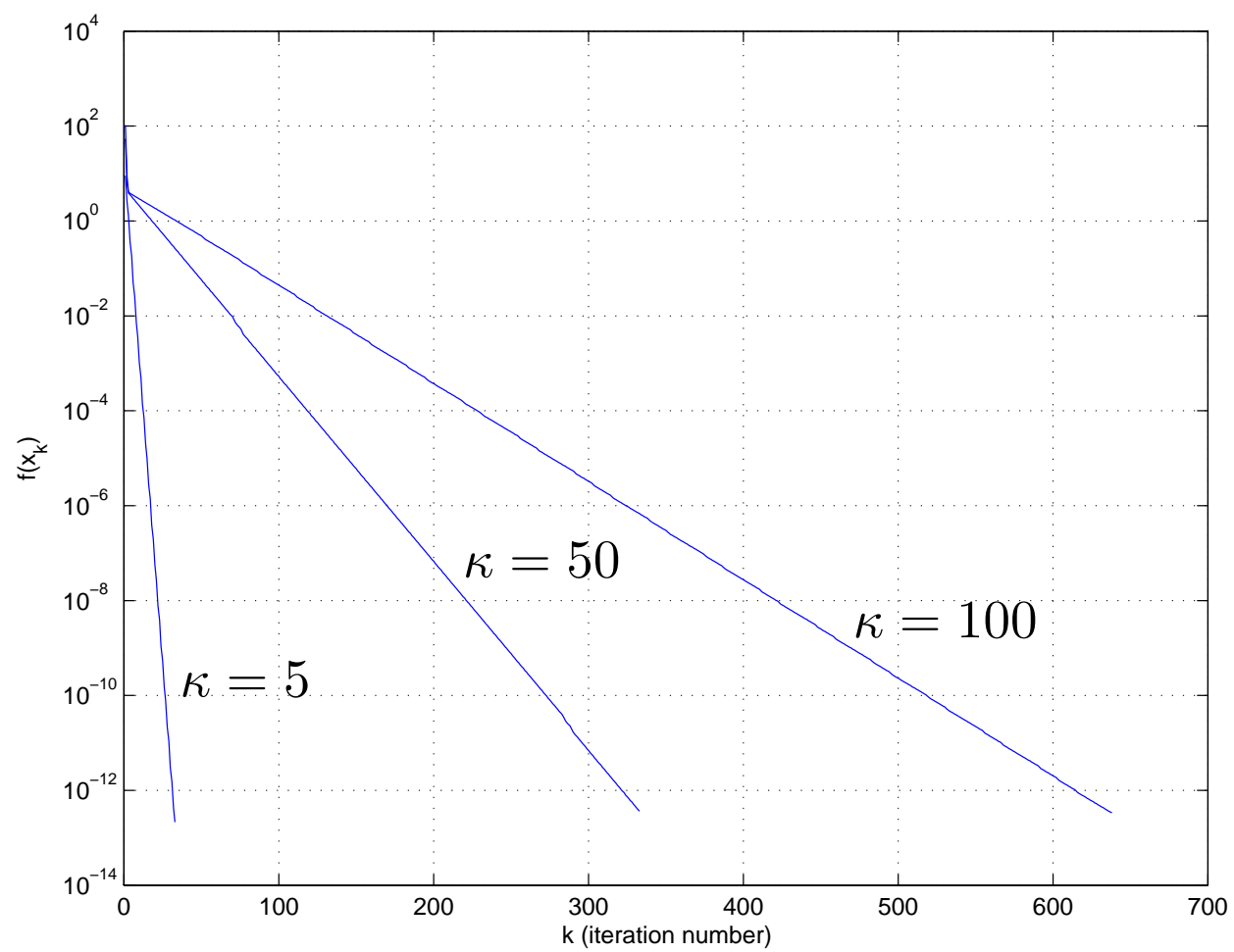
$\kappa = 50$  (number of iterations: 332)





$\kappa = 100$  (number of iterations: 637)

Sequence  $\{f(x_k)\}_{k \geq 0}$



**Definition (rates of convergence)** Let  $x_k \rightarrow x^*$  ( $x_k, x^* \in \mathbb{R}^n$ ) with  $x_k \neq x^*$ . We say

- $x_k \rightarrow x^*$  linearly if

$$\exists_{0 < r < 1} : \frac{\|x_{k+1} - x^*\|}{\|x_k - x^*\|} \leq r \text{ for all } k \text{ sufficiently large}$$

- $x_k \rightarrow x^*$  superlinearly if

$$\frac{\|x_{k+1} - x^*\|}{\|x_k - x^*\|} \rightarrow 0$$

- $x_k \rightarrow x^*$  quadratically if

$$\exists_{r > 0} : \frac{\|x_{k+1} - x^*\|}{\|x_k - x^*\|^2} \leq r \text{ for all } k \text{ sufficiently large}$$

Examples:

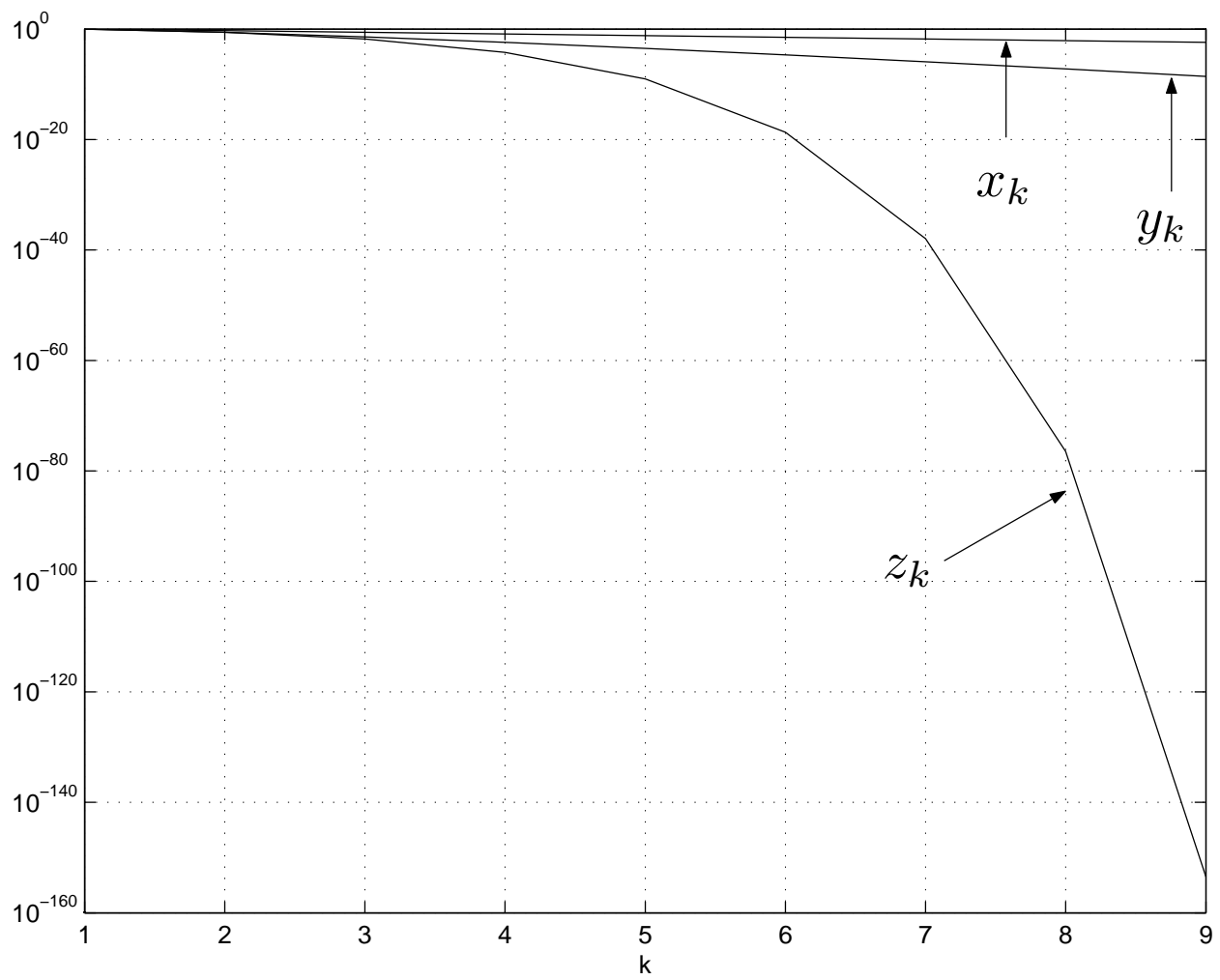
$$x_k = \left(\frac{1}{2}\right)^{k-1} \rightarrow 0 \text{ linearly}$$

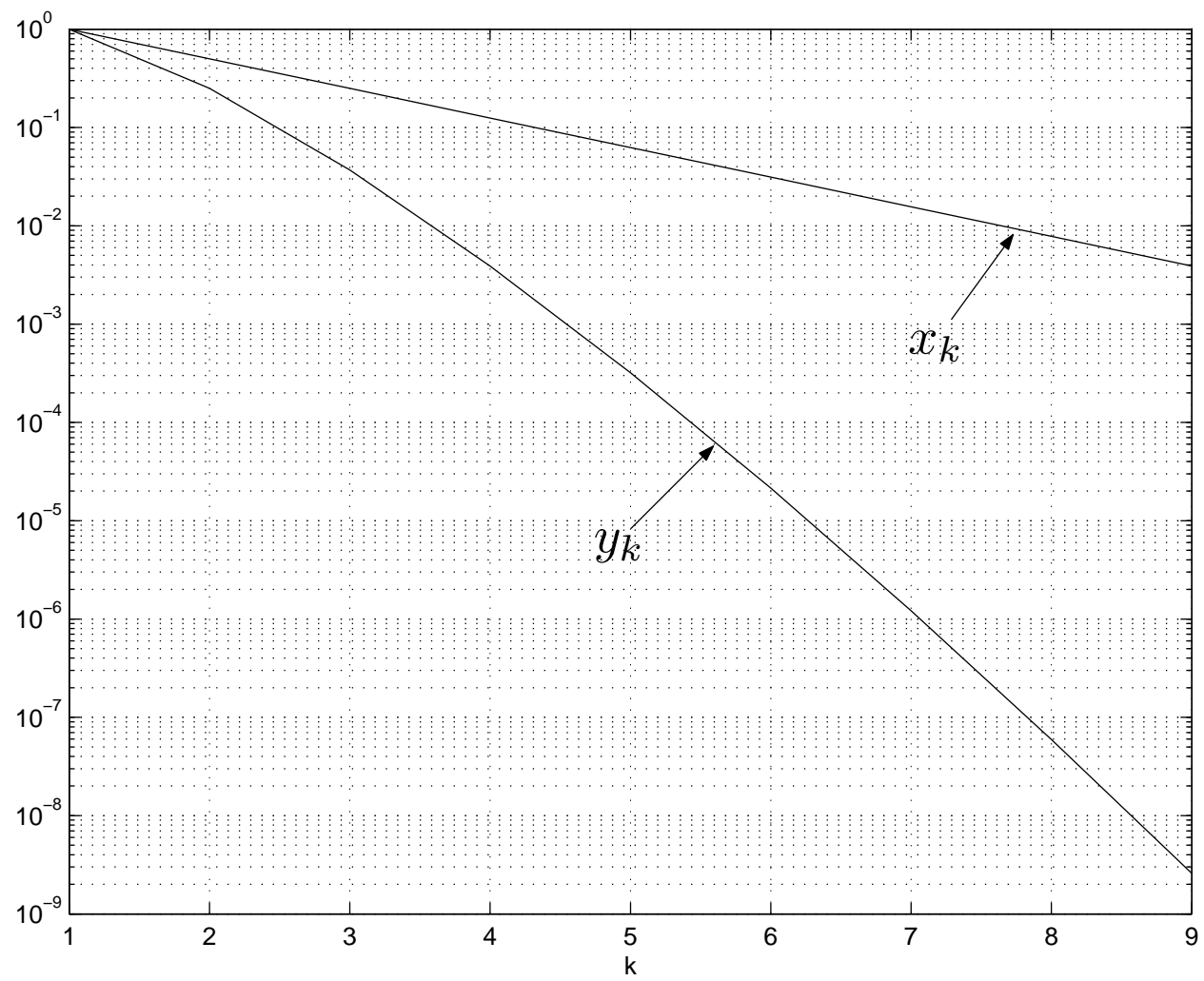
$$y_k = \frac{1}{k^k} \rightarrow 0 \text{ superlinearly}$$

$$z_k = 4 \left(\frac{1}{2}\right)^{2^k} \rightarrow 0 \text{ quadratically}$$

First few terms of the sequences...

$k$	$x_k$	$y_k$	$z_k$
1	1	1	1
2	$5.0 \cdot 10^{-1}$	$2.5 \cdot 10^{-1}$	$2.5 \cdot 10^{-1}$
3	$2.5 \cdot 10^{-1}$	$3.7 \cdot 10^{-2}$	$1.6 \cdot 10^{-2}$
4	$1.3 \cdot 10^{-1}$	$3.9 \cdot 10^{-3}$	$6.1 \cdot 10^{-5}$
5	$6.3 \cdot 10^{-2}$	$3.2 \cdot 10^{-4}$	$9.3 \cdot 10^{-10}$
6	$3.1 \cdot 10^{-2}$	$2.1 \cdot 10^{-5}$	$2.2 \cdot 10^{-19}$
7	$1.6 \cdot 10^{-2}$	$1.2 \cdot 10^{-6}$	$1.2 \cdot 10^{-38}$
8	$7.8 \cdot 10^{-3}$	$6.0 \cdot 10^{-8}$	$3.5 \cdot 10^{-77}$
9	$3.9 \cdot 10^{-3}$	$2.6 \cdot 10^{-9}$	$3.0 \cdot 10^{-154}$





Under mild assumptions, the steepest descent method  $d_k = -\nabla f(x_k)$  yields linear convergence



Newton algorithm (first motivation):

- $x_k$  is the current iterate
- the 2nd order model of  $f$  centered at  $x_k$  is

$$\hat{f}_k(x) = f(x_k) + \nabla f(x_k)^\top (x - x_k) + \frac{1}{2}(x - x_k)^\top \nabla^2 f(x_k)(x - x_k)$$

- minimize 2nd order model to obtain the next iterate  $x_{k+1}$ :

$$x_{k+1} = \arg \min_x \hat{f}_k(x) = x_k + d_k \quad \text{where} \quad d_k = - [\nabla^2 f(x_k)]^{-1} \nabla f(x_k)$$

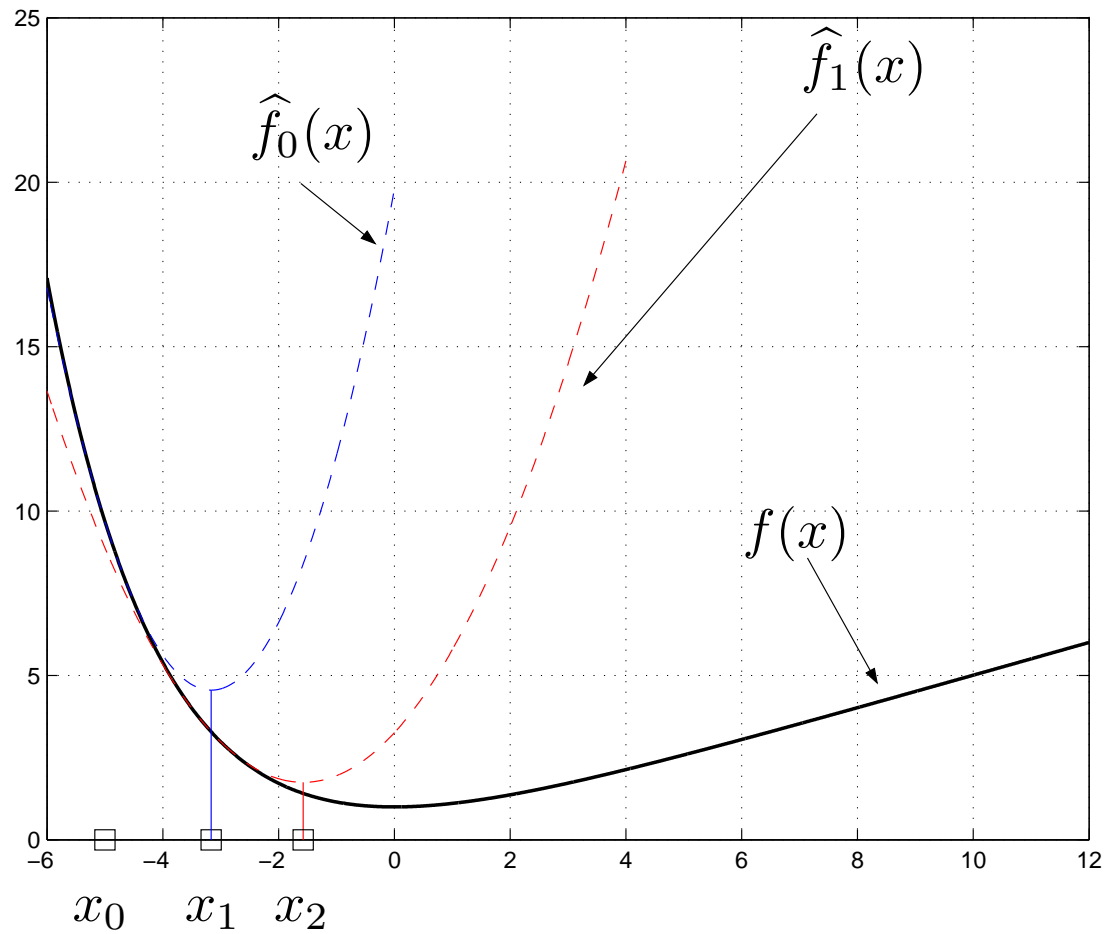
Note: it is assumed that  $\nabla^2 f(x_k)$  is nonsingular

The algorithm

$$x_{k+1} = x_k + d_k \quad \text{with} \quad d_k = - [\nabla^2 f(x_k)]^{-1} \nabla f(x_k)$$

is called the pure Newton algorithm

Example:  $f : \mathbb{R} \rightarrow \mathbb{R}$ ,  $f(x) = e^{-\frac{1}{2}x} + \frac{1}{2}x$



Newton algorithm (second motivation):

- Newton method for solving a system of nonlinear equations

$$\left\{ \begin{array}{l} F_1(x_1, \dots, x_n) = 0 \\ F_2(x_1, \dots, x_n) = 0 \\ \vdots \\ F_n(x_1, \dots, x_n) = 0 \end{array} \right. \Leftrightarrow F(x) = 0$$

is

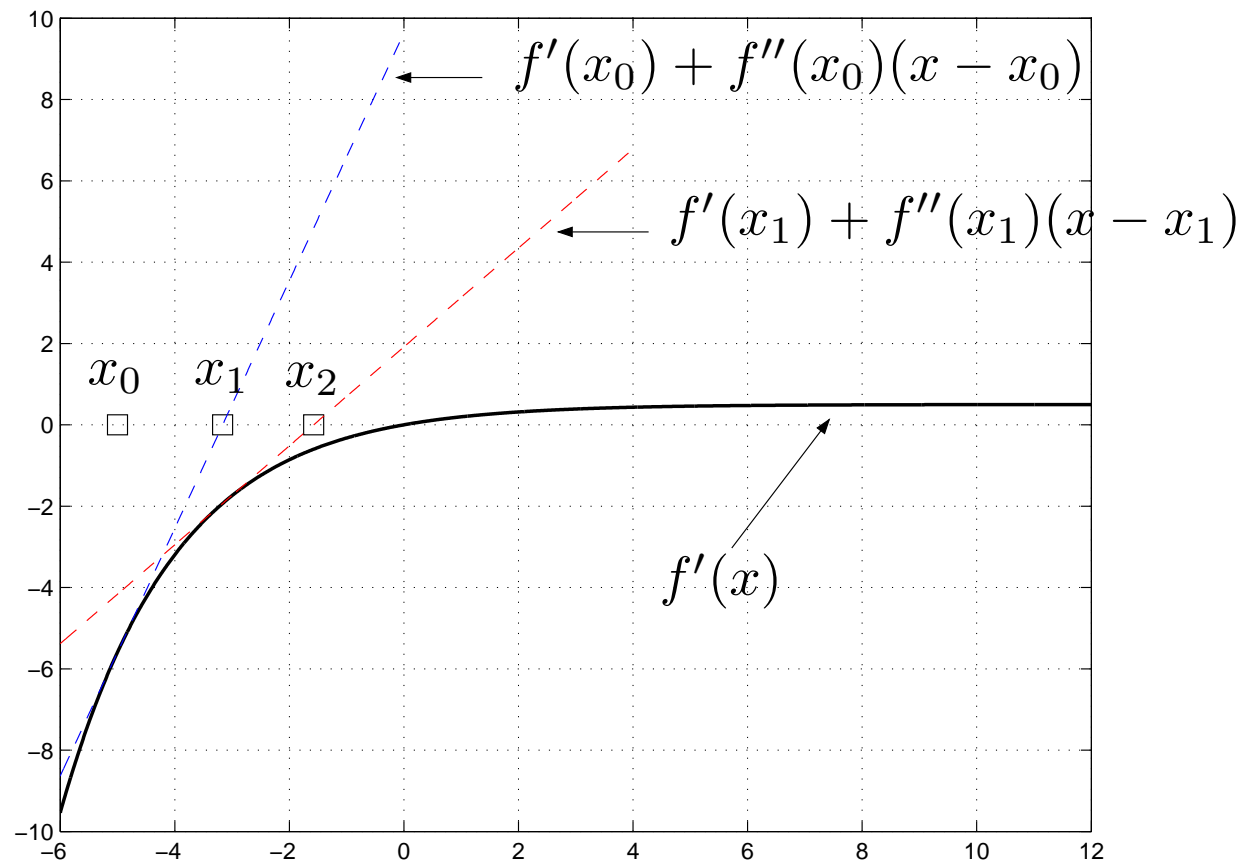
$$x_{k+1} = x_k - DF(x_k)^{-1} F(x_k).$$

Note: the inverse  $DF(x_k)^{-1}$  is supposed to exist

- the pure Newton algorithm corresponds to the special case  $\nabla f(x) = 0$ :

$$x_{k+1} = x_k - [\nabla^2 f(x_k)]^{-1} \nabla f(x_k)$$

Example:  $f : \mathbb{R} \rightarrow \mathbb{R}$ ,  $f(x) = e^{-\frac{1}{2}x} + \frac{1}{2}x$



- The pure Newton algorithm  $x_{k+1} = x_k + d_k$  is convergent only near a solution
- The Newton direction

$$d_k = - [\nabla^2 f(x_k)]^{-1} \nabla f(x_k)$$

is a descent direction for  $f$  at  $x_k$ :

$$\langle d_k, \nabla f(x_k) \rangle = -\nabla f(x_k)^\top [\nabla^2 f(x_k)]^{-1} \nabla f(x_k) < 0$$

- Use a line-search to make it globally convergent:  $x_{k+1} = x_k + t_k d_k$

- The quantity

$$\lambda(x) := \sqrt{\nabla f(x)^\top [\nabla^2 f(x)]^{-1} \nabla f(x)} = \sqrt{-g_k^\top d_k}$$

is called the Newton decrement at  $x$

- Property of the Newton decrement:

$$\frac{1}{2}\lambda^2(x) = f(x) - \min \left\{ f(x) + \nabla f(x)^\top (y - x) + \frac{1}{2} (y - x)^\top \nabla^2 f(x) (y - x) : y \right\}$$

- Affine invariant stopping criterion:

$$\lambda^2(x) < \epsilon$$

- The Newton algorithm solves a sequence of unconstrained quadratic problems

1.      **initialization**    ►    choose  $x_0 \in \text{dom } f$  and tolerance  $\epsilon > 0$
2.                            ►    choose  $0 < \alpha < 0.5$  and  $0 < \beta < 1$  (Armijo)
3.                            ►    set  $k = 0$
4.      **loop**                    ►    compute  $g_k = \nabla f(x_k)$
5.                            ►    compute  $H_k = \nabla^2 f(x_k)$
6.                            ►    solve the linear system  $H_k d_k = -g_k$
7.                            ►    if  $g_k^\top H_k^{-1} g_k = -g_k^\top d_k < \epsilon$  stop
8.                            ►     $t = 1$
9.                            ►    while  $f(x_k + t d_k) > f(x_k) + t \alpha g_k^\top d_k$   
   $t \leftarrow \beta t$
10.                           ►     $x_{k+1} = x_k + t d_k$
11.                           ►     $k \leftarrow k + 1$  and return to loop



**Theorem (Convergence of Newton method)** Let  $x_0 \in \text{dom } f$  and suppose there exists  $m > 0$  such that

$$\nabla^2 f(x) \succeq mI \quad \text{for all } x \in S_{f(x_0)} f.$$

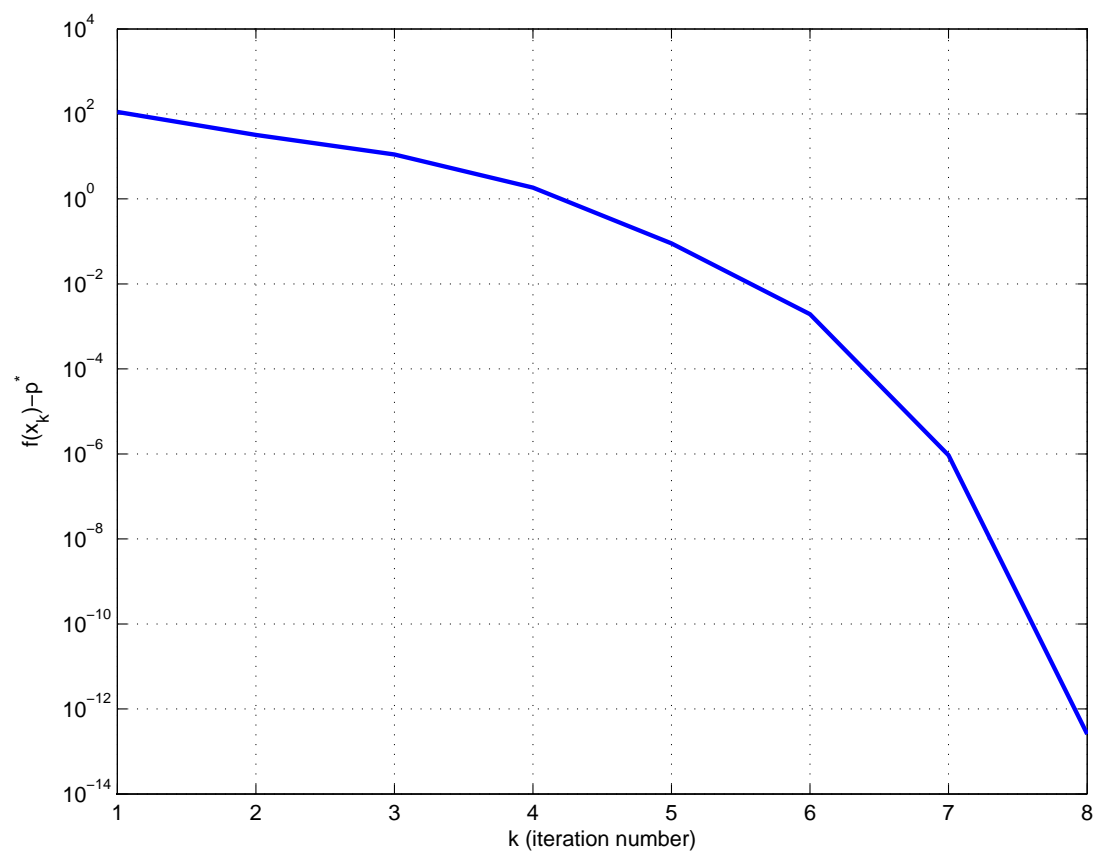
Also assume that  $\nabla^2 f(x)$  is Lipschitz continuous on  $S_{f(x_0)} f$ . Then,

$$x_k \rightarrow x^* \quad \text{quadratically.}$$

Example:

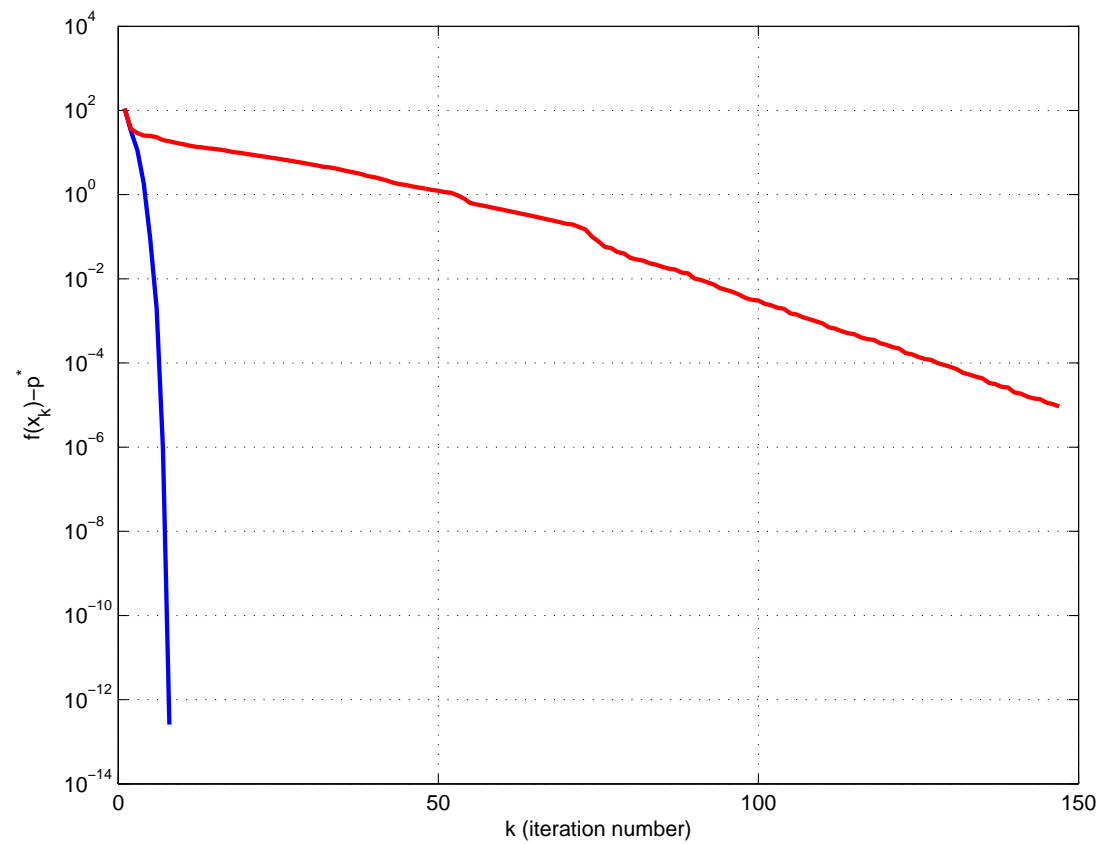
$$f(x) = c^\top x - \sum_{i=1}^m \log(b_i - a_i^\top x)$$

## Newton method



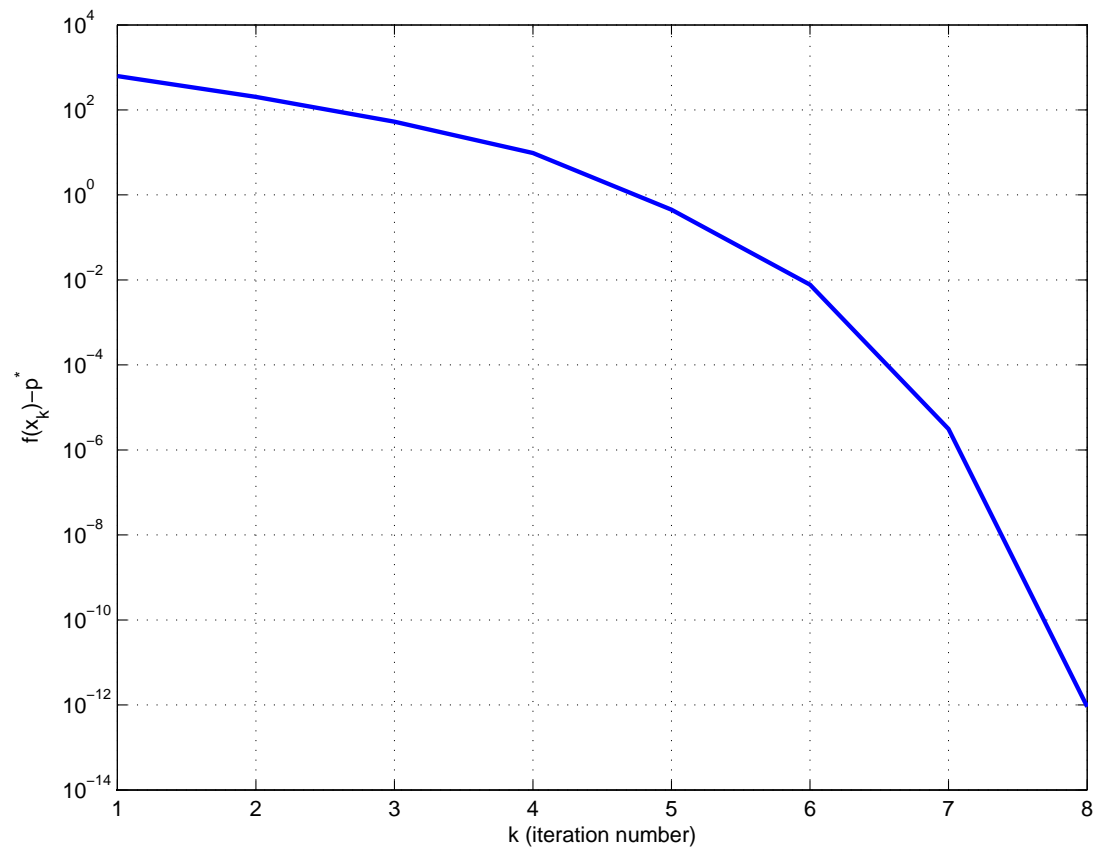
$$n = 8, m = 50, \kappa(\nabla^2 f(x^*)) \simeq 42$$

## Newton and gradient descent methods



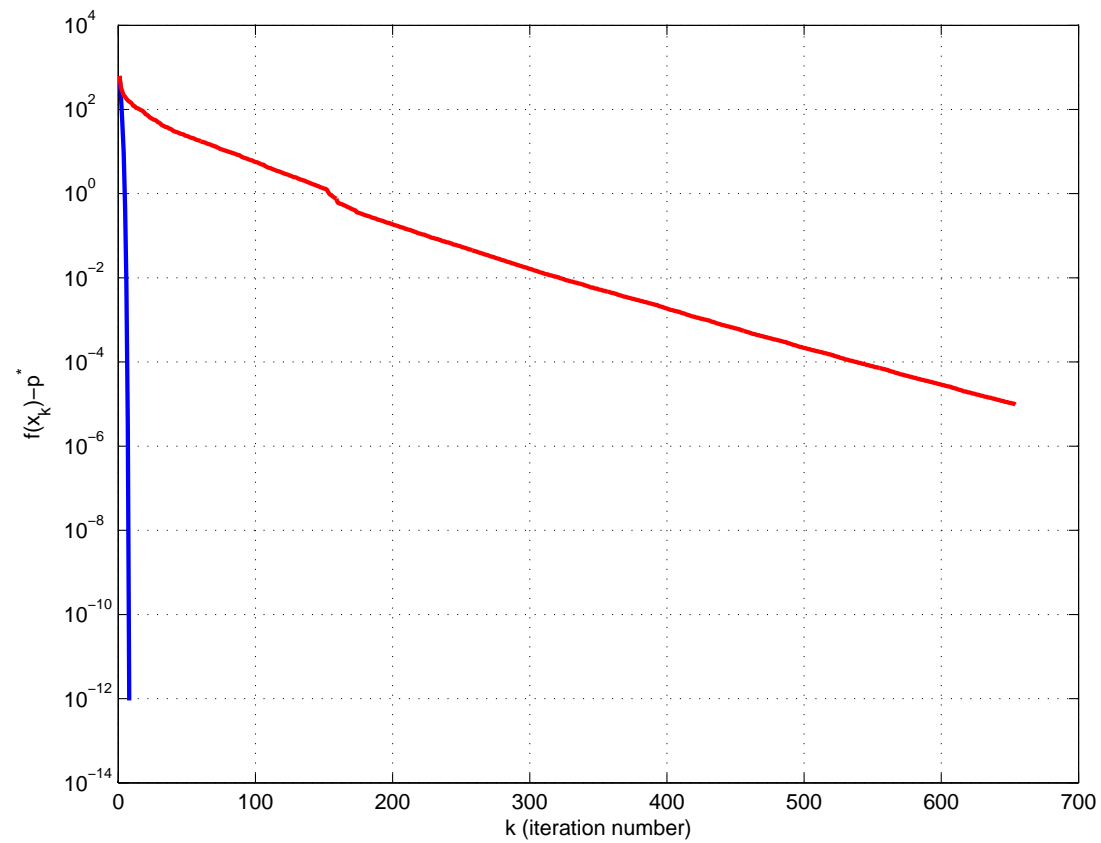
$$n = 8, m = 50, \kappa(\nabla^2 f(x^*)) \simeq 42$$

## Newton method



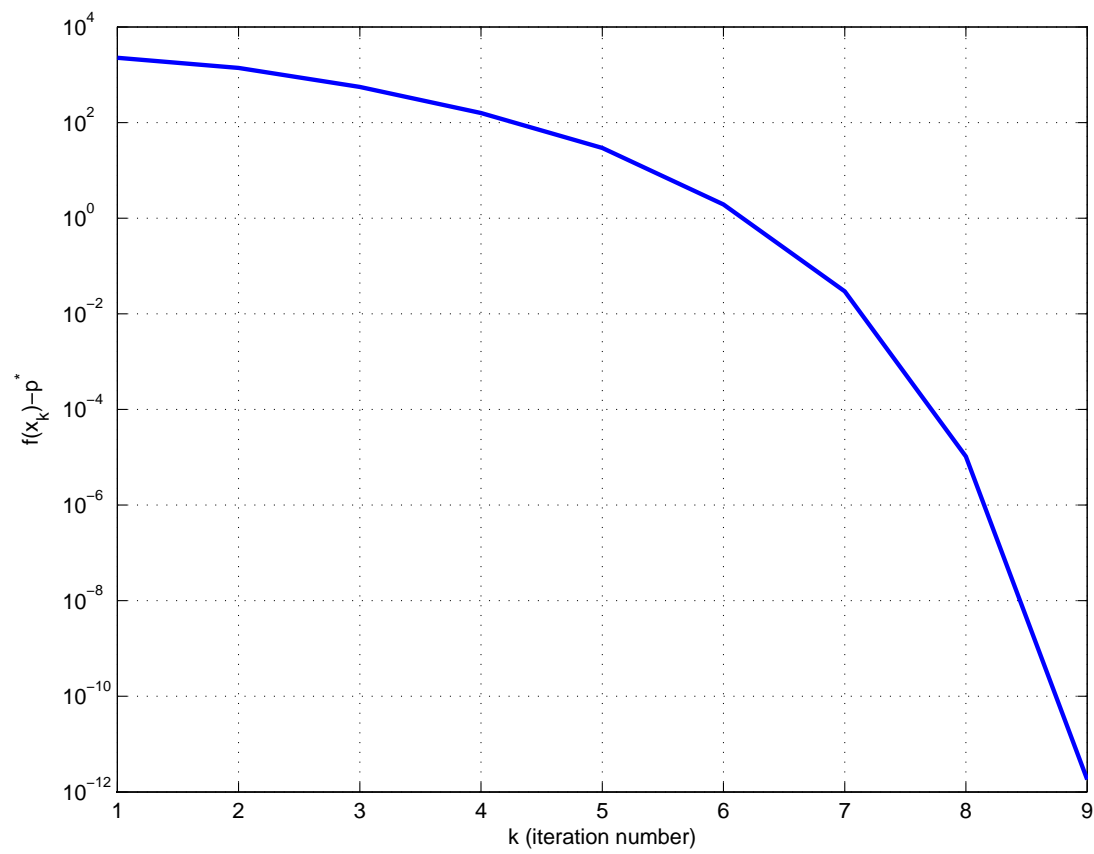
$$n = 50, m = 250, \kappa(\nabla^2 f(x^*)) \simeq 253$$

## Newton and gradient descent methods



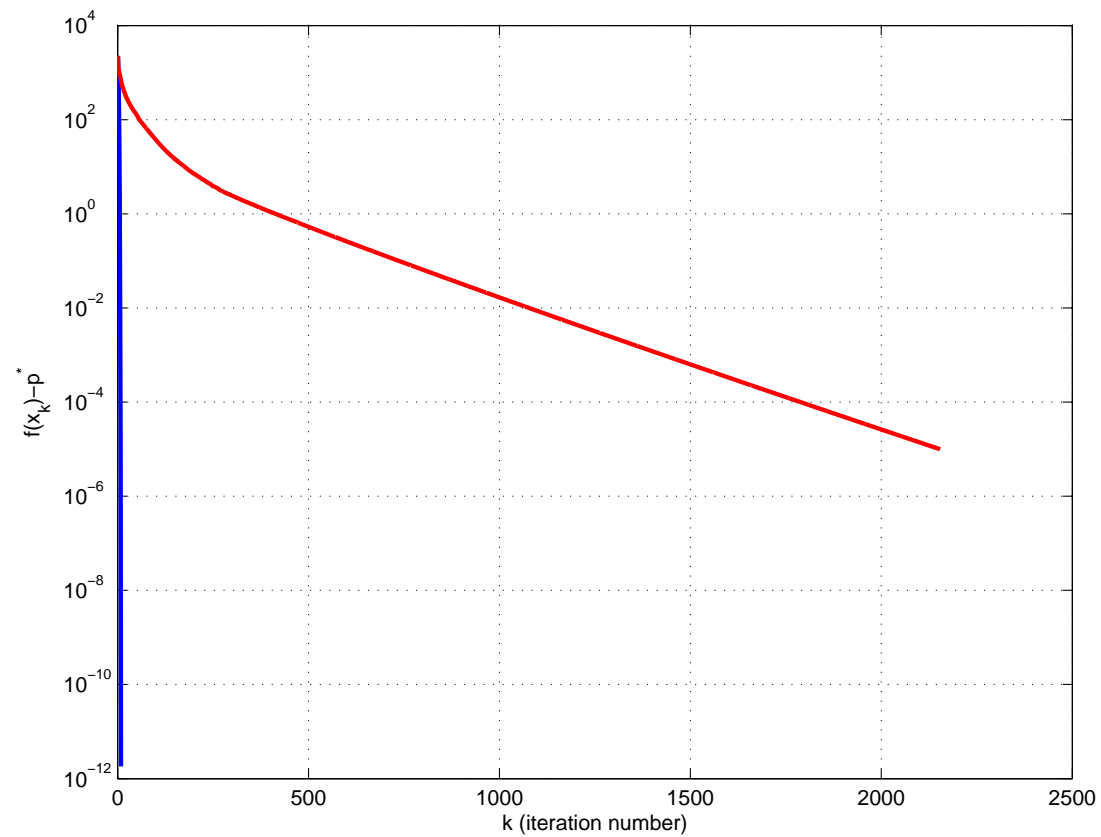
$$n = 50, m = 250, \kappa(\nabla^2 f(x^*)) \simeq 253$$

## Newton method



$$n = 200, m = 800, \kappa(\nabla^2 f(x^*)) \simeq 777$$

## Newton and gradient descent methods

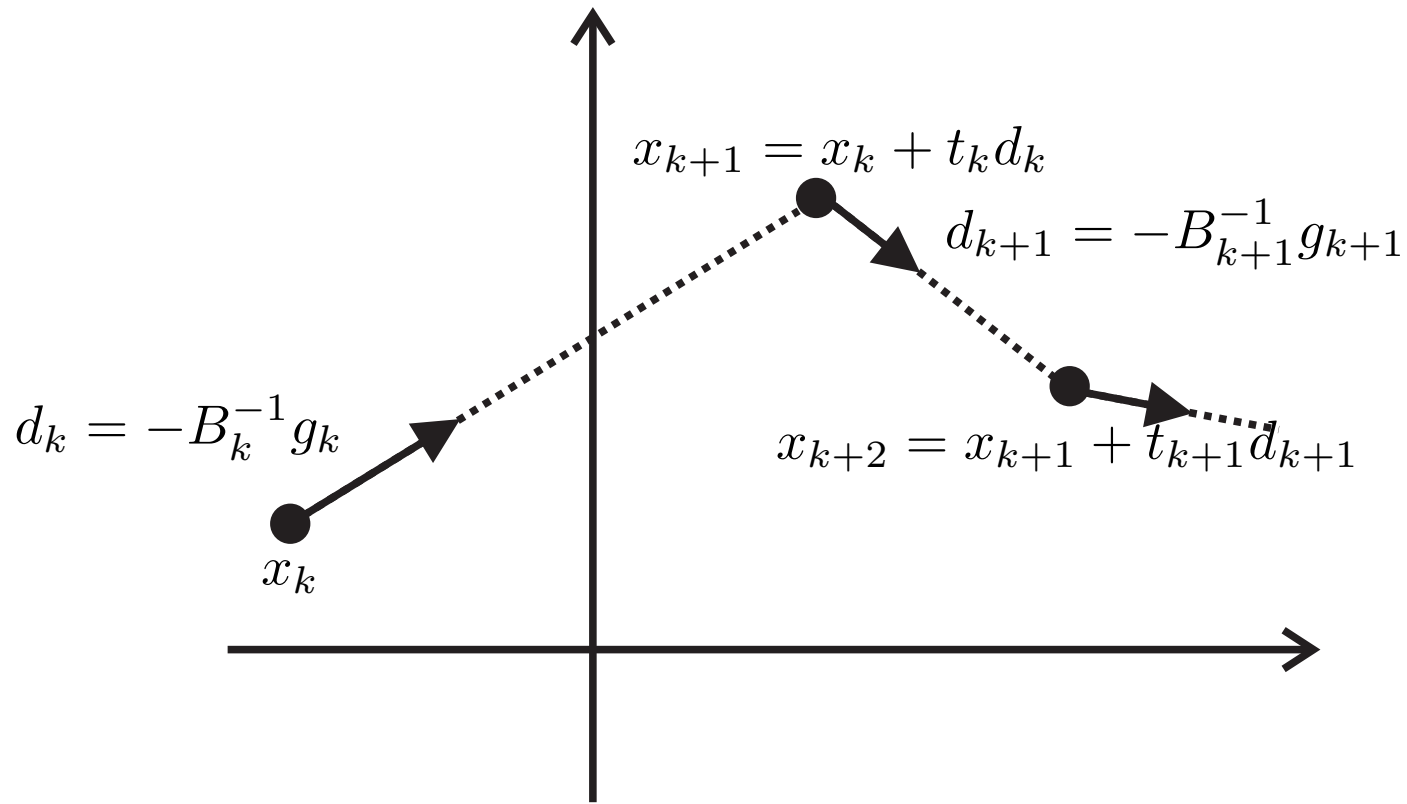


$$n = 200, m = 800, \kappa(\nabla^2 f(x^*)) \simeq 777$$



## Quasi-Newton BFGS algorithm

- BFGS = Broyden, Fletcher, Goldfarb and Shanno
- Newton method needs  $H_k = \nabla^2 f(x_k)$  at each iteration
- idea: instead, use an approximation  $B_k$



- **Initialization:**  $B_0 = I_n$  or a better guess
- **Update:** how to update  $B_k \rightarrow B_{k+1}$  ?

Quadratic model  $m(x)$  for  $f(x)$  centered at  $x_{k+1}$ :

$$m(x) = f(x_{k+1}) + \nabla f(x_{k+1})^\top (x - x_{k+1}) + \frac{1}{2} (x - x_{k+1})^\top B_{k+1} (x - x_{k+1})$$

“Nice” properties that  $B_{k+1}$  should possess:

- $\nabla m(x_k) = g_k$  which leads to the secant equation

$$B_{k+1} \underbrace{(x_{k+1} - x_k)}_{s_k} = \underbrace{(g_{k+1} - g_k)}_{y_k}$$

- $B_{k+1} \succ 0$  to ensure that  $d_{k+1} = -B_{k+1}^{-1}g_k$  is a descent direction
- $B_{k+1} \simeq B_k$  ( $f$  is  $C^2$  implies  $\nabla^2 f(x)$  changes continuously with  $x$ )

First two properties might conflict: for given  $s_k, y_k$  there exist no

$$B_{k+1} \succ 0 \text{ such that } B_{k+1}s_k = y_k$$

Example:  $s_k = 1, y_k = -1$  implies  $B_{k+1} = y_k/s_k = -1$

**Fact.** Let  $s, y \in \mathbb{R}^n$  be given with  $s \neq 0$ . Then,

$$\exists B \succ 0 : y = Bs \quad \Leftrightarrow \quad y^\top s > 0.$$

Proof ( $\Leftarrow$ ):  $B = I_n - \frac{ss^\top}{s^\top s} + \frac{yy^\top}{y^\top s}$  works.

Note: if  $\nabla^2 f(x) \succ 0$  for all  $x$  then  $y_k^\top s_k > 0$

- Translate desired properties into an optimization problem:

$$\begin{aligned} &\text{minimize} && \phi(B) := \text{tr} \left( B_k^{-1/2} B B_k^{-1/2} \right) - \log \det \left( B_k^{-1/2} B B_k^{-1/2} \right) - n \\ &\text{subject to} && B s_k = y_k \\ &&& B \succ 0 \end{aligned}$$

- $\phi(B)$  is a “distance” measure between  $B$  and  $B_k$ :

$$\phi(B) \geq 0 \text{ with equality if and only if } B = B_k$$

- Solution is

$$B_{k+1} = B_k - \frac{B_k s_k s_k^\top B_k}{s_k^\top B_k s_k} + \frac{y_k y_k^\top}{y_k^\top s_k}.$$

This rank-2 update is known as the **BFGS update**

We can propagate directly the inverse  $W_{k+1} = B_{k+1}^{-1}$ :

$$W_{k+1} = (I - \rho_k s_k y_k^\top) W_k (I - \rho_k y_k s_k^\top) + \rho_k s_k s_k^\top \quad \text{where} \quad \rho_k = \frac{1}{y_k^\top s_k}$$

1.       **initialization**       ▶   choose  $x_0 \in \text{dom } f$ ,  $W_0 \succ 0$  and  $\epsilon > 0$
2.                               ▶   choose  $0 < \alpha < 0.5$  and  $0 < \beta < 1$  (Armijo)
3.                               ▶   set  $k = 0$
4.       **loop**               ▶   compute  $g_k = \nabla f(x_k)$
5.                               ▶   if  $\|g_k\| < \epsilon$  stop
6.                               ▶   compute search direction  $d_k = -W_k g_k$
7.                               ▶    $t = 1$
8.                               ▶   while  $f(x_k + td_k) > f(x_k) + t\alpha g_k^\top d_k$   
   $t \leftarrow \beta t$
9.                               ▶    $x_{k+1} = x_k + td_k$
10.                              ▶    $W_{k+1} = (I - \rho_k s_k y_k^\top) W_k (I - \rho_k y_k s_k^\top) +$   
   $\rho_k s_k s_k^\top$  with  $\rho_k = 1/y_k^\top s_k$
11.                              ▶    $k \leftarrow k + 1$  and return to loop

**Theorem (Convergence of Quasi-Newton BFGS method)** Let  $x_0 \in \text{dom } f$  and suppose there exists  $m > 0$  such that

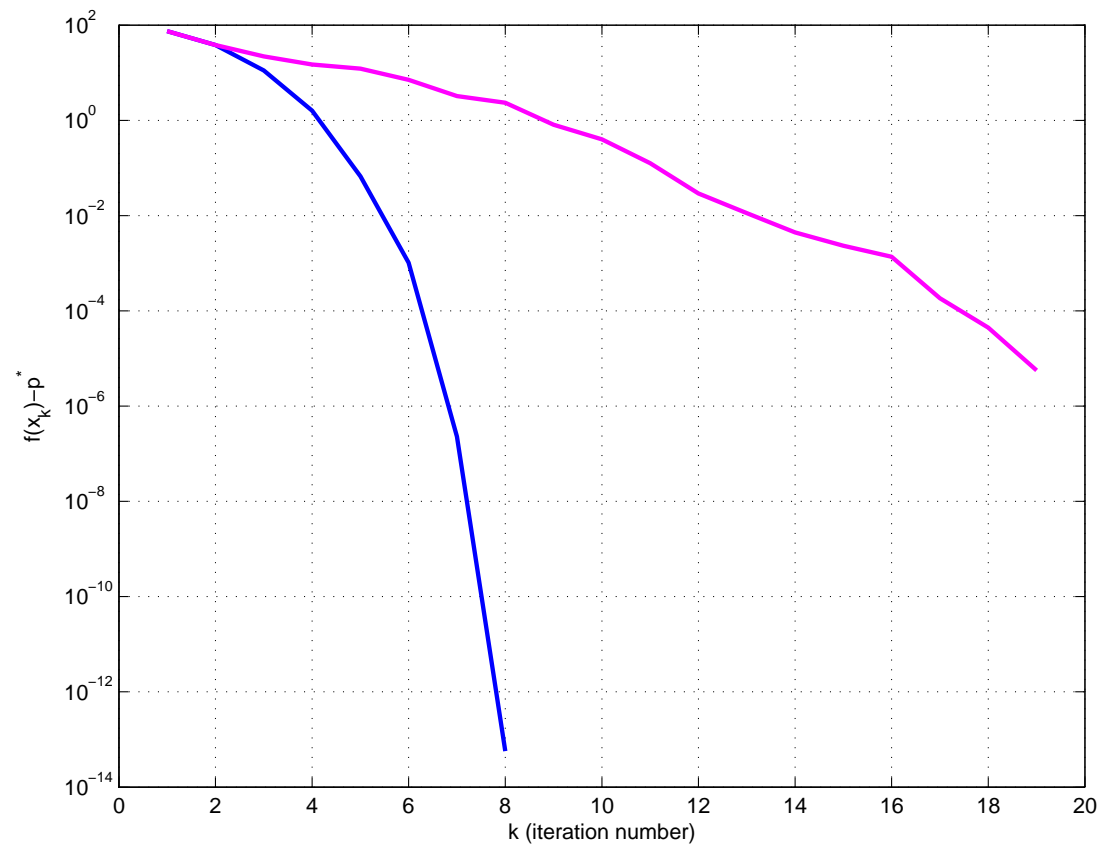
$$\nabla^2 f(x) \succeq mI \quad \text{for all } x \in S_{f(x_0)} f.$$

Also assume that  $\nabla^2 f(x)$  is Lipschitz continuous on  $S_{f(x_0)} f$ . Then,

$$x_k \rightarrow x^* \quad \text{superlinearly.}$$

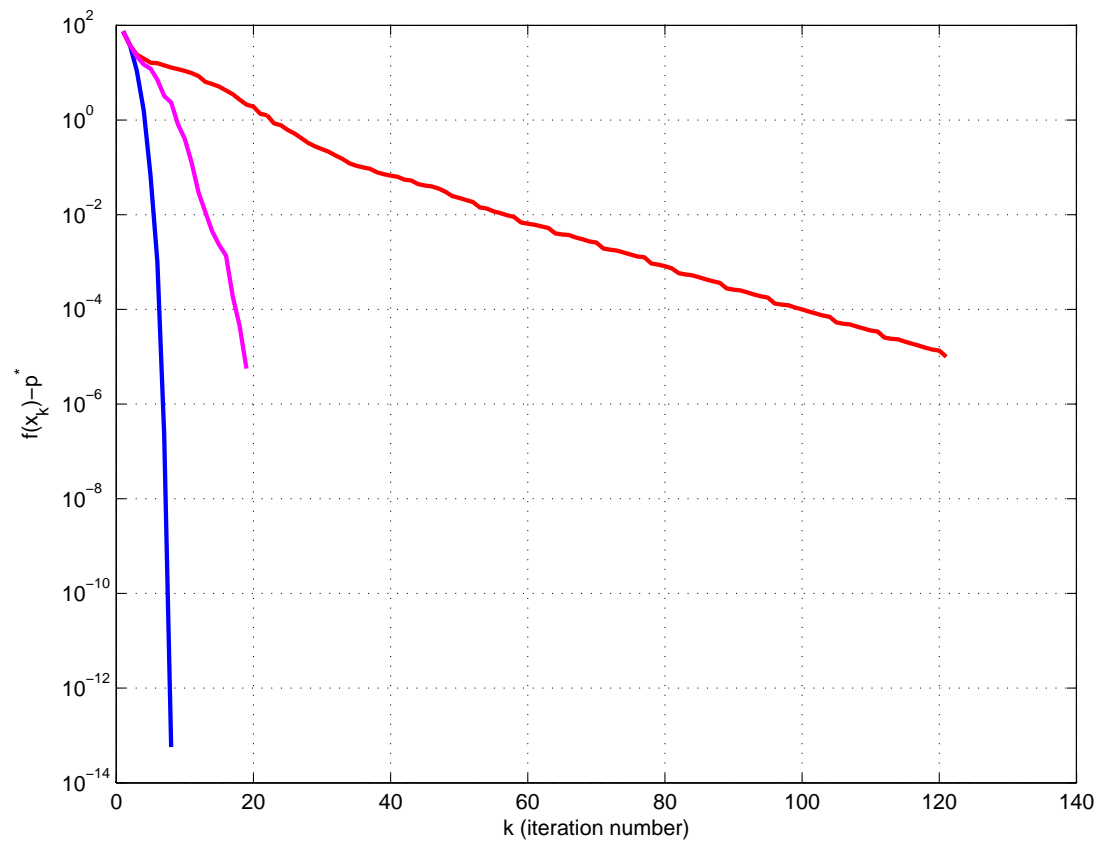


## Newton and Quasi-Newton BFGS methods



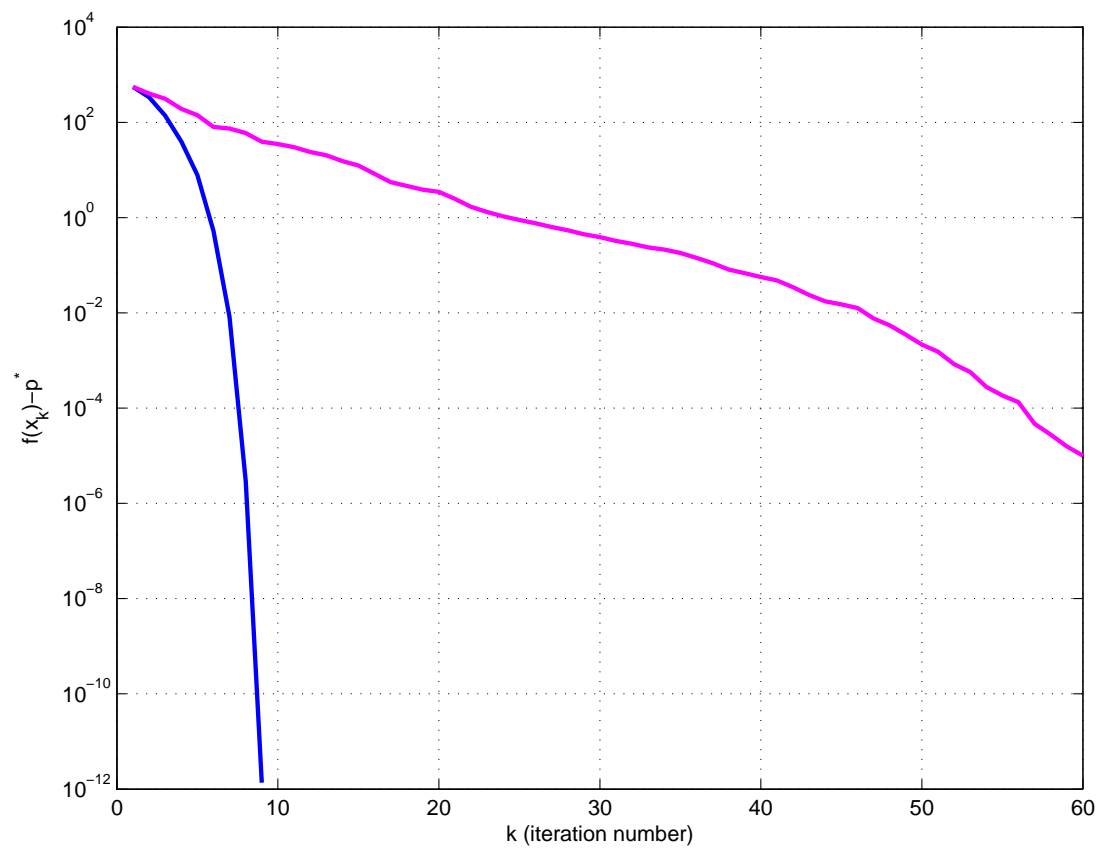
$$n = 8, m = 50, \kappa(\nabla^2 f(x^*)) \simeq 786$$

## Newton, Quasi-Newton BFGS and gradient descent methods



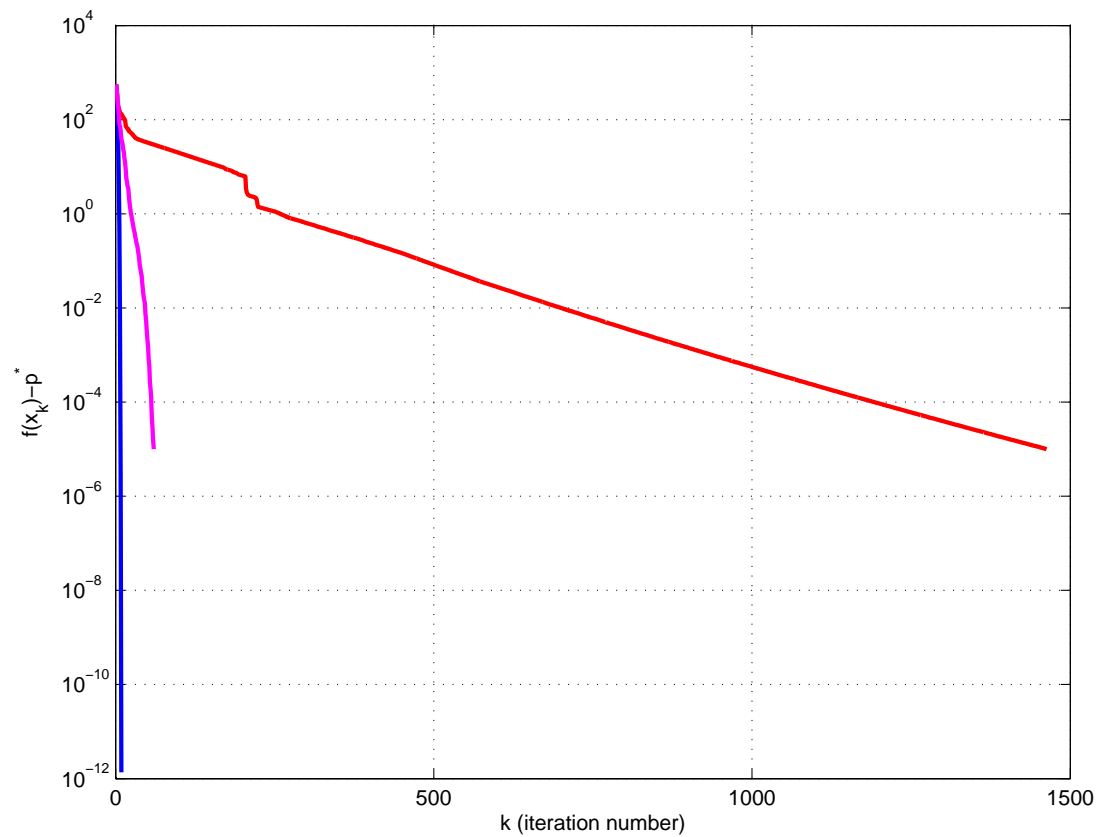
$$n = 8, m = 50, \kappa(\nabla^2 f(x^*)) \simeq 786$$

## Newton and Quasi-Newton BFGS methods



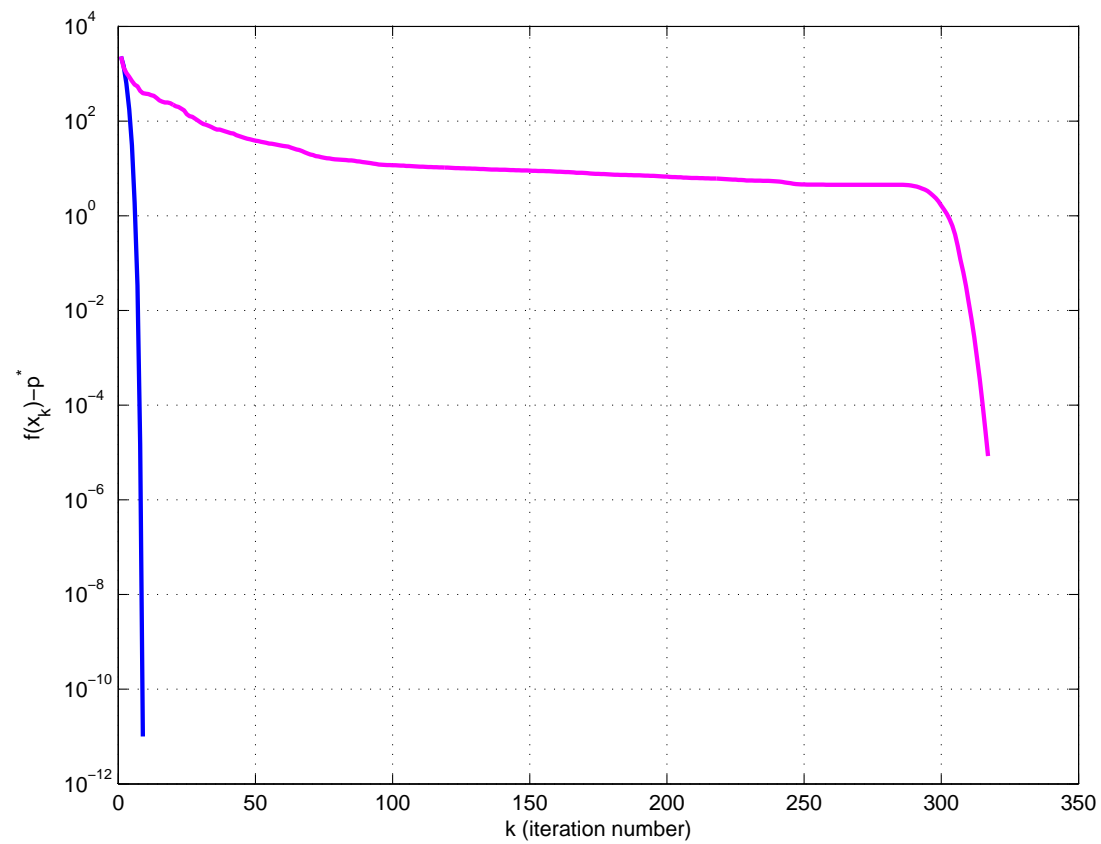
$$n = 50, m = 250, \kappa(\nabla^2 f(x^*)) \simeq 851$$

## Newton, Quasi-Newton BFGS and gradient descent methods



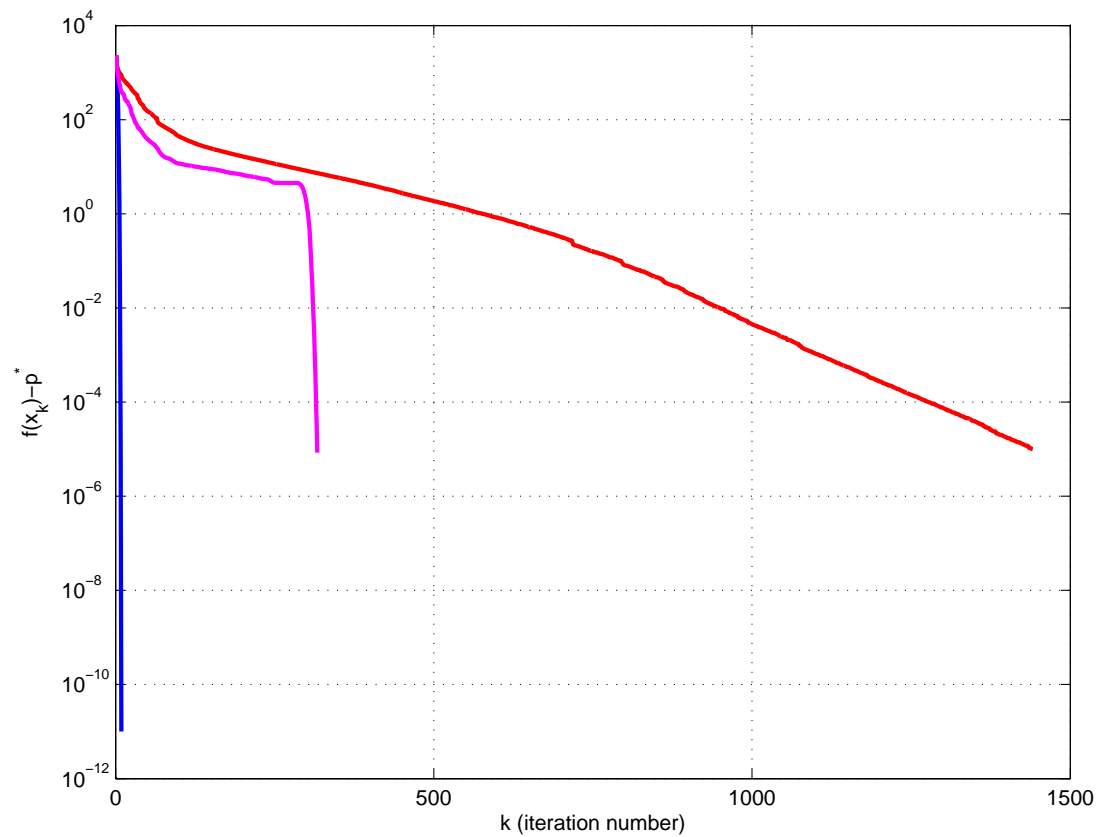
$$n = 50, m = 250, \kappa(\nabla^2 f(x^*)) \simeq 851$$

## Newton and Quasi-Newton BFGS methods



$$n = 200, m = 800, \kappa(\nabla^2 f(x^*)) \simeq 873$$

## Newton, Quasi-Newton BFGS and gradient descent methods



$$n = 200, m = 800, \kappa(\nabla^2 f(x^*)) \simeq 873$$

# **Constrained minimization (equality constraints)**

Optimization problem:

$$\begin{array}{ll} \text{minimize} & f(x) \\ \text{subject to} & Ax = b \end{array}$$

Assumptions:

- $f : \mathbb{R}^n \rightarrow \mathbb{R} \cup \{+\infty\}$  is convex
- $\text{dom } f = \{x : f(x) < +\infty\}$  is open
- $f$  is closed
- $f$  is twice continuously differentiable
- problem is solvable:  $\exists x^* : f(x^*) = p^* := \inf\{f(x) : Ax = b\}$
- matrix  $A$  has full row rank (otherwise, drop redundant eqs)



Two main approaches:

- eliminate equality constraints and switch to unconstrained problem
- maintain equality constraints

Eliminate equality constraints:

- parameterize constraint set as

$$\{x : Ax = b\} = \{x_0 + By : y\}$$

where  $Ax_0 = b$  and  $B$  contains a basis for  $\text{Ker } A$

- solve equivalent unconstrained optimization problem

$$y^* = \arg \min_y f(x_0 + By)$$

- recover solution

$$x^* = x_0 + By^*$$

- sparsity structure in  $A$  might be lost

Maintain equality constraints (first motivation):

- $x_k$  is the current iterate
- the 2nd order model of  $f$  centered at  $x_k$  is

$$\hat{f}_k(x) = f(x_k) + \nabla f(x_k)^\top (x - x_k) + \frac{1}{2}(x - x_k)^\top \nabla^2 f(x_k)(x - x_k)$$

- minimize 2nd order model to obtain the next iterate  $x_{k+1}$ :

$$\begin{aligned} x_{k+1} = \arg \min \quad & \hat{f}_k(x) \\ \text{subject to} \quad & Ax = b \end{aligned} \quad = x_k + d_k$$

where

$$\begin{bmatrix} \nabla^2 f(x_k) & A^\top \\ A & 0 \end{bmatrix} \begin{bmatrix} d_k \\ \lambda_k \end{bmatrix} = - \begin{bmatrix} \nabla f(x_k) \\ 0 \end{bmatrix}$$

Maintain equality constraints (second motivation):

- KKT system is

$$\begin{cases} \nabla f(x) + A^\top \lambda = 0 \\ Ax - b = 0 \end{cases} \Leftrightarrow r(x, \lambda) = 0$$

- linearize around the current iterate  $x_k$ :

$$r(x_k + d_k, \lambda_k) \simeq (\nabla f(x_k) + \nabla^2 f(x_k) d_k + A^\top \lambda_k, A(x_k + d_k) - b)$$

- solve for  $(d_k, \lambda_k)$ :

$$\begin{bmatrix} \nabla^2 f(x_k) & A^\top \\ A & 0 \end{bmatrix} \begin{bmatrix} d_k \\ \lambda_k \end{bmatrix} = - \begin{bmatrix} \nabla f(x_k) \\ 0 \end{bmatrix}$$

- $x_{k+1} = x_k + d_k$

1.       **initialization**   ▶   choose  $x_0 \in \text{dom } f$  ( $Ax_0 = b$ ) and  $\epsilon > 0$
2.                           ▶   choose  $0 < \alpha < 0.5$  and  $0 < \beta < 1$  (Armijo)
3.                           ▶   set  $k = 0$
4.       **loop**           ▶   compute  $g_k = \nabla f(x_k)$
5.                           ▶   compute  $H_k = \nabla^2 f(x_k)$
6.                           ▶   solve  $\begin{bmatrix} H_k & A^\top \\ A & 0 \end{bmatrix} \begin{bmatrix} d_k \\ \lambda_k \end{bmatrix} = \begin{bmatrix} -g_k \\ 0 \end{bmatrix}$
7.                           ▶   if  $-g_k^\top d_k < \epsilon$  **stop**
8.                           ▶    $t = 1$
9.                           ▶   **while**  $f(x_k + td_k) > f(x_k) + t\alpha g_k^\top d_k$   
   $t \leftarrow \beta t$
10.                          ▶    $x_{k+1} = x_k + td_k$
11.                          ▶    $k \leftarrow k + 1$  and return to loop

# **Constrained minimization (inequality constraints)**

Optimization problem:

$$\begin{array}{ll}\text{minimize} & f(x) \\ \text{subject to} & Ax = b \\ & g_i(x) \leq 0 \quad i = 1, \dots, m\end{array}$$

Assumptions:

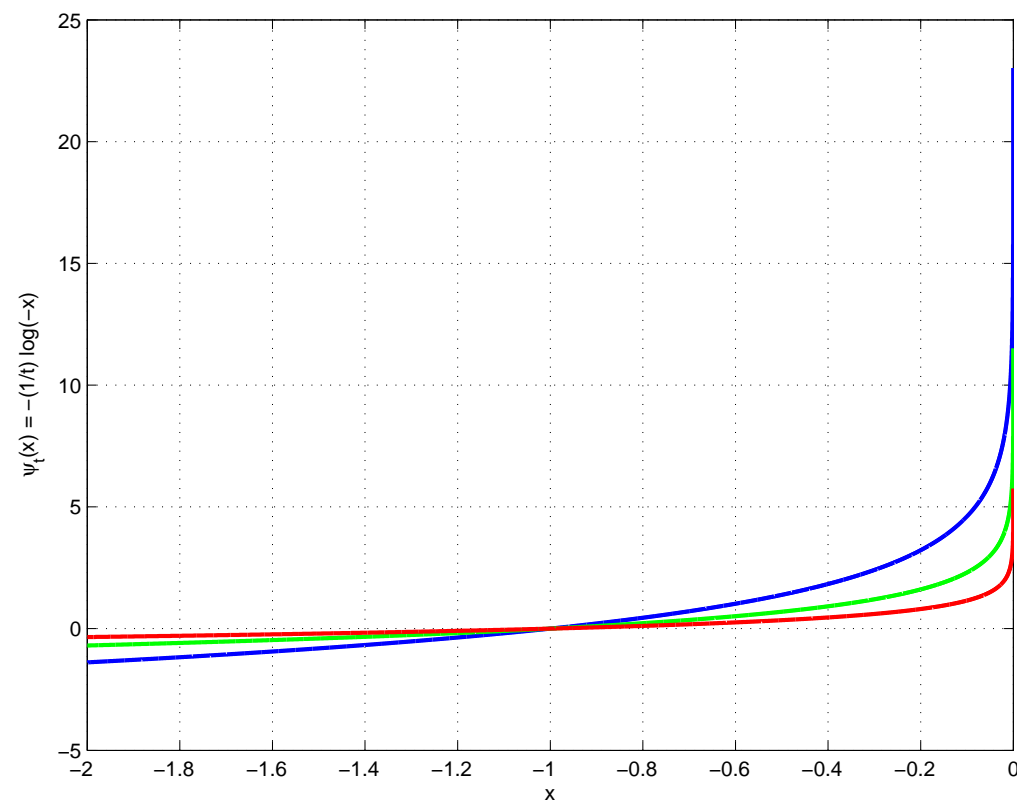
- $f, g_i : \mathbb{R}^n \rightarrow \mathbb{R} \cup \{+\infty\}$  are convex
- $\text{dom } f, \text{dom } g_i$  are open
- $f, g_i$  are closed and twice continuously differentiable
- problem is solvable:

$$\exists x^* : f(x^*) = p^* := \inf\{f(x) : Ax = b, g(x) = (g_1(x), \dots, g_m(x)) \leq 0\}$$

- matrix  $A$  has full row rank
- there is a Slater point:  $\exists_{x \in \text{dom } f} : Ax = b, g(x) < 0$

Logarithmic barrier:

$$\psi_t : \mathbb{R} \rightarrow \mathbb{R} \cup \{+\infty\} \quad \psi_t(x) = -\frac{1}{t} \log(-x) \quad \text{dom } \psi_t = -\mathbb{R}_{++}$$



Plot of  $\psi_t$  for  $t = 0.5$ ,  $t = 1$  and  $t = 2$



- Consider the approximation

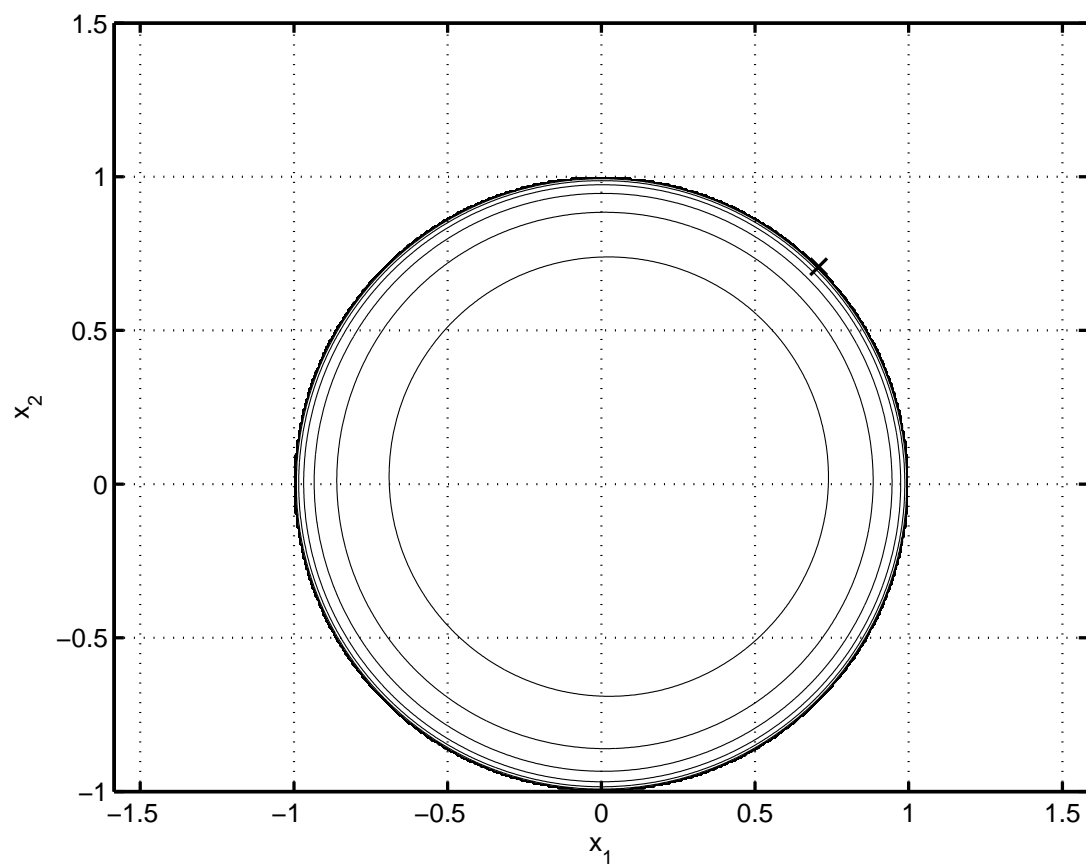
$$\begin{array}{ll}\text{minimize} & f(x) + \sum_{i=1}^m -(1/t) \log(-g_i(x)) \\ \text{subject to} & Ax = b\end{array}$$

- Example:

$$\begin{array}{ll}\text{minimize} & -x_1 - x_2 \\ \text{subject to} & x_1^2 + x_2^2 \leq 1\end{array}$$

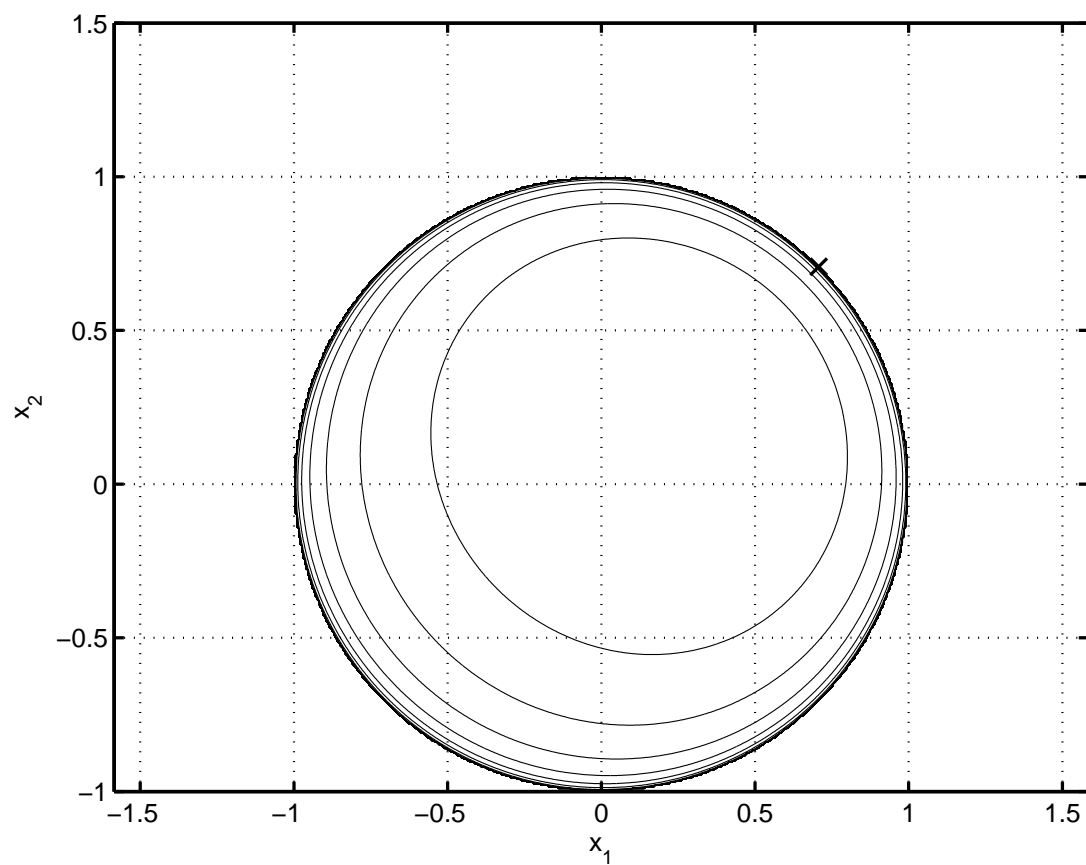
Solution is  $(x_1^*, x_2^*) = (\frac{1}{\sqrt{2}}, \frac{1}{\sqrt{2}})$

Level sets for  $\phi_t(x_1, x_2) = -x_1 - x_2 - (1/t) \log(1 - x_1^2 - x_2^2)$



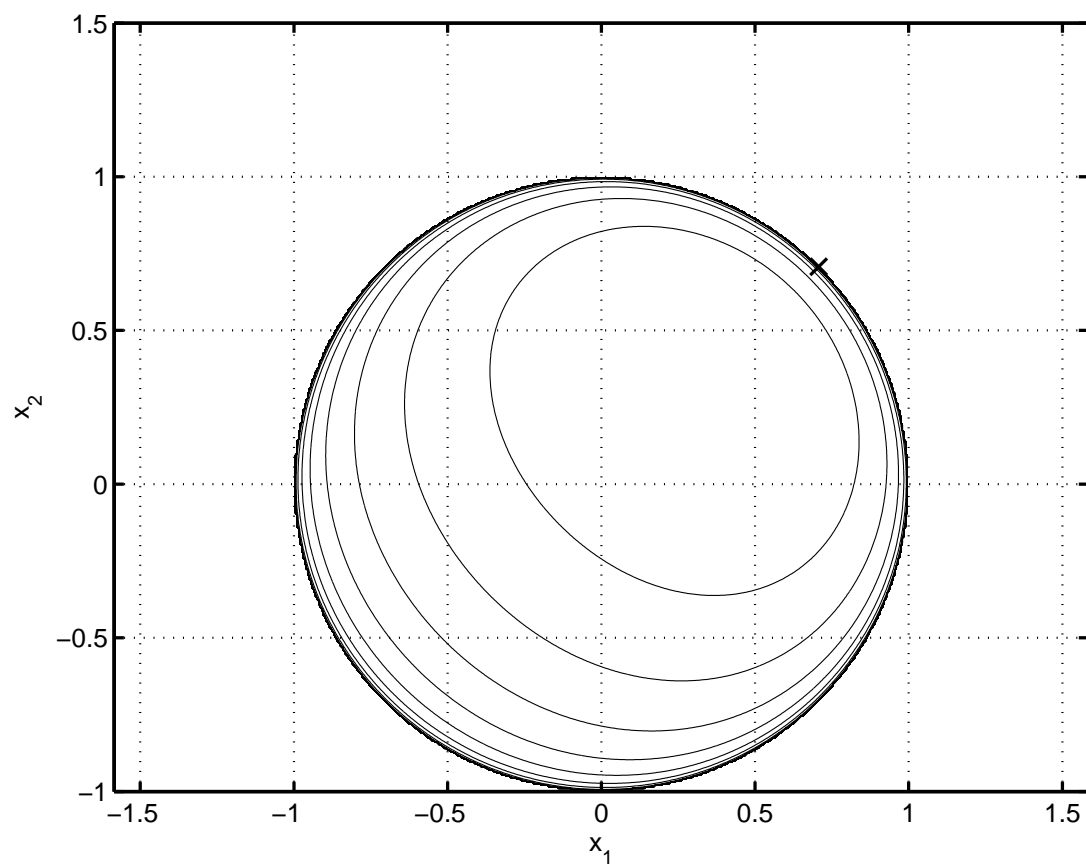
$t = 0.1$

Level sets for  $\phi_t(x_1, x_2) = -x_1 - x_2 - (1/t) \log(1 - x_1^2 - x_2^2)$



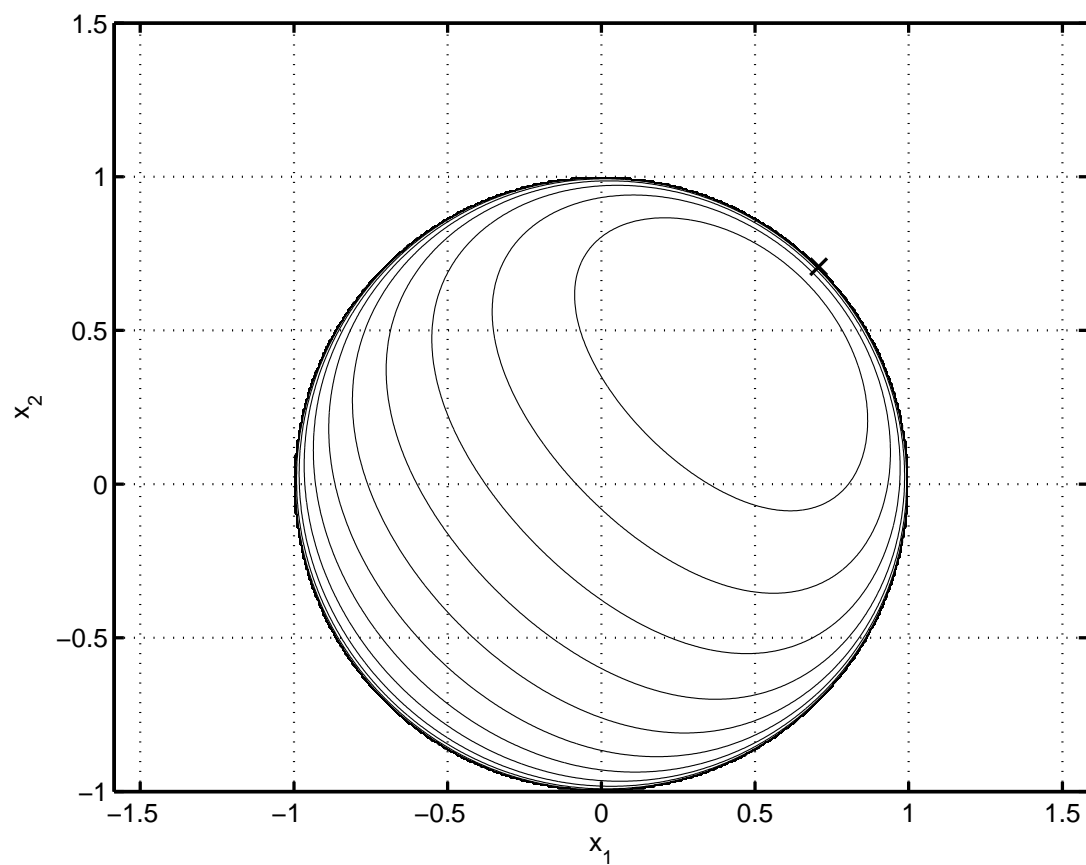
$t = 0.5$

Level sets for  $\phi_t(x_1, x_2) = -x_1 - x_2 - (1/t) \log(1 - x_1^2 - x_2^2)$



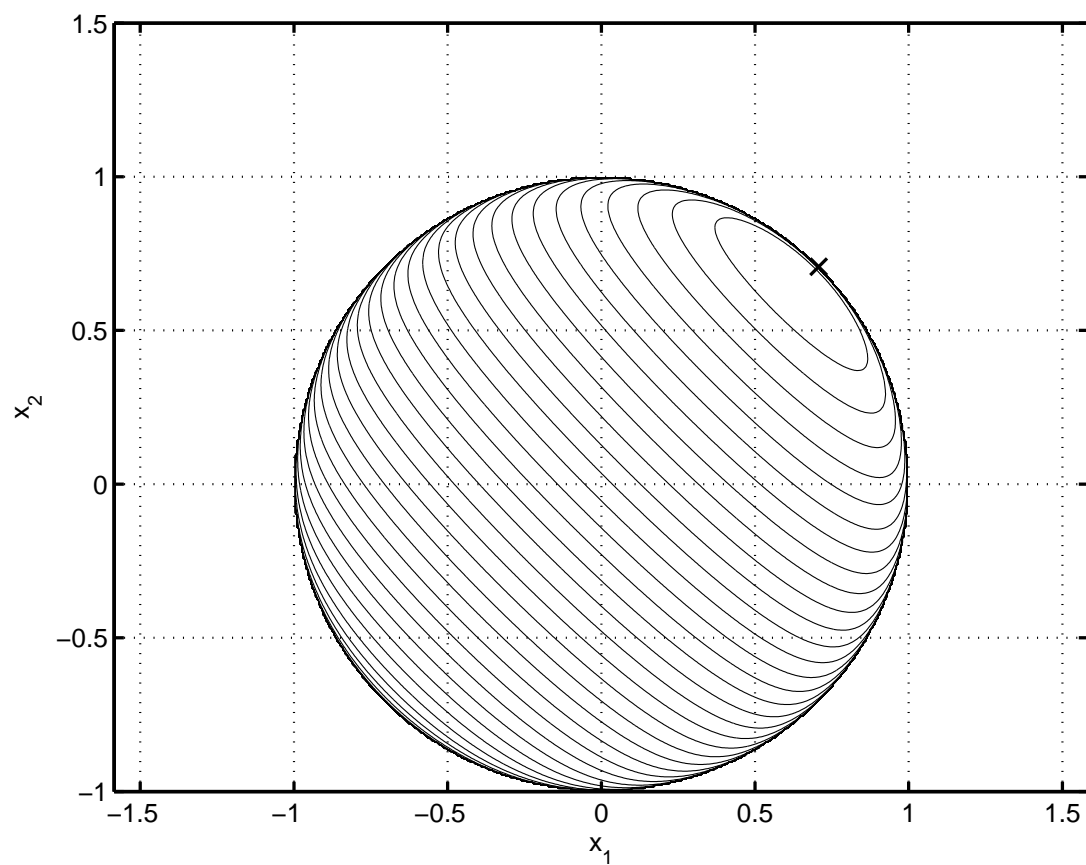
$t = 1$

Level sets for  $\phi_t(x_1, x_2) = -x_1 - x_2 - (1/t) \log(1 - x_1^2 - x_2^2)$



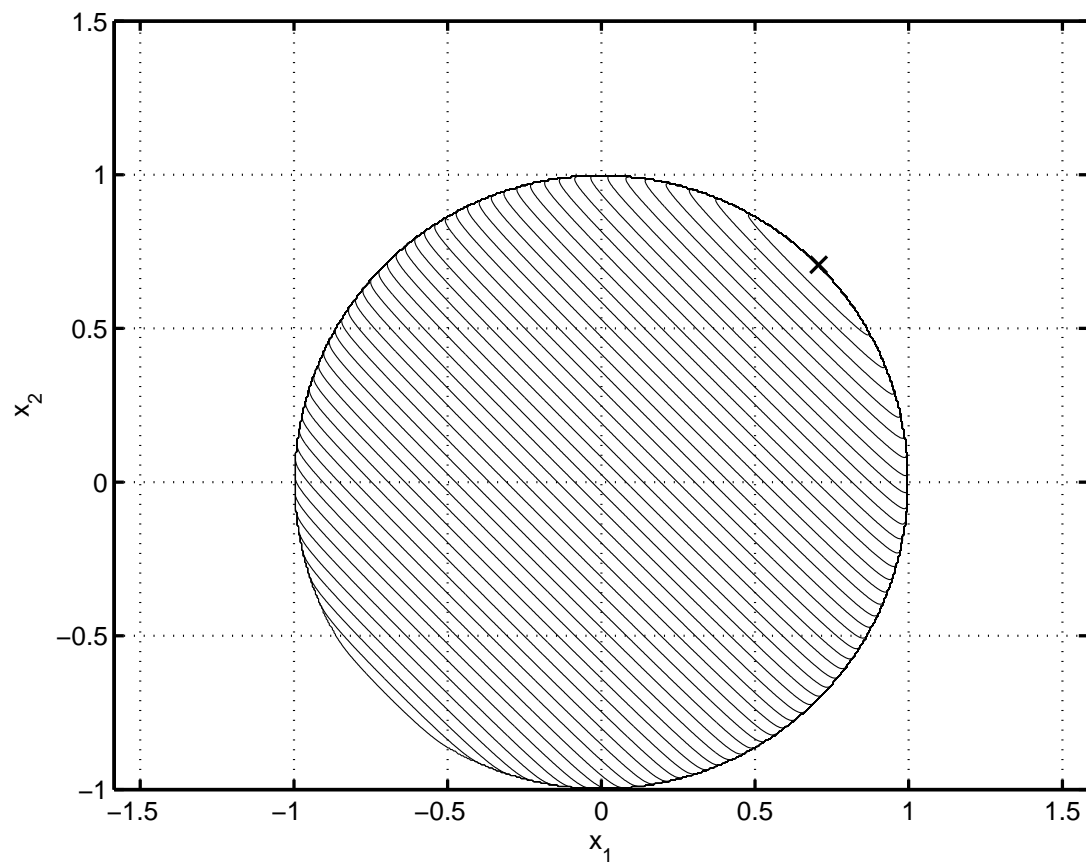
$t = 2$

Level sets for  $\phi_t(x_1, x_2) = -x_1 - x_2 - (1/t) \log(1 - x_1^2 - x_2^2)$



$t = 10$

Level sets for  $\phi_t(x_1, x_2) = -x_1 - x_2 - (1/t) \log(1 - x_1^2 - x_2^2)$



$t = 100$

- Assume an unique minimizer  $x^*(t)$  for

$$\begin{array}{ll} \text{minimize} & f(x) + \sum_{i=1}^m -(1/t) \log(-g_i(x)) \\ \text{subject to} & Ax = b \end{array}$$

- $x^*(t)$  is (at most)  $m/t$  suboptimal:

$$f(x^*(t)) - p^* \leq \frac{m}{t}$$

- Central path =  $\{x^*(t) : t > 0\}$
- Following the central path leads to a solution



1.     **initialization**     ▶      $t_0 > 0$ , strictly feasible  $x_0$ ,  $\mu > 1$ ,  $\epsilon > 0$
2.                             ▶     initialize  $k = 0$
3.     **loop**                 ▶     solve  $\min_{Ax=b} t_k f(x) + \sum_{i=1}^m -\log(-g_i(x))$   
   starting from  $x_k$
4.                             ▶      $x_{k+1} = x^*(t_k)$
5.                             ▶     if  $m/t_k < \epsilon$  then stop
6.                             ▶      $t_{k+1} = \mu t_k$
7.                             ▶      $k \leftarrow k + 1$  and return to loop

- To construct a strictly feasible point  $x_0$  use a “Phase I” problem:

minimize  $s$

subject to  $Ax = b$

$$g_i(x) \leq s \quad i = 1, \dots, m$$

- Can start from  $x \in \bigcap_{i=1}^m \text{dom } f_i$  with  $Ax = b$  and

$$s > \max\{g_i(x) : i = 1, \dots, m\}$$

Optimization problem:

$$\begin{array}{ll}\text{minimize} & f(x) \\ \text{subject to} & Ax = b \\ & g_i(x) \preceq_{K_i} 0 \quad i = 1, \dots, m\end{array}$$

Assumptions:

- $K_i$  are proper cones
- $f$  is convex and  $g_i$  is  $K_i$ -convex
- $\text{dom } f, \text{dom } g_i$  are open and  $f, g_i$  are twice continuously differentiable
- problem is solvable:

$$\exists x^* : f(x^*) = p^* := \inf\{f(x) : Ax = b, g_i(x) \preceq_{K_i} 0, i = 1, \dots, m\}$$

- matrix  $A$  has full row rank
- there is a Slater point:  
 $\exists x \in \text{dom } f : Ax = b, g_i(x) \prec_{K_i} 0, i = 1, \dots, m$

A generalized logarithm for a proper cone  $K \subset \mathbb{R}^n$  is a function

$$\psi : \mathbb{R}^n \rightarrow \mathbb{R} \cup \{+\infty\} \quad \text{dom } \psi = \text{int } K$$

such that

- $\psi$  is closed, concave, twice continuously differentiable with

$$\nabla^2 \psi(x) \prec 0 \quad \text{for all } x \succ_K 0$$

- there is  $\theta > 0$  (degree of  $\psi$ ) such that

$$\psi(tx) = \psi(x) + \theta \log(t) \quad \text{for all } x \succ_K 0 \text{ and } t > 0$$

Examples:

- nonnegative orthant:  $K = \mathbb{R}_+^n$

$$\psi(x_1, \dots, x_n) = \sum_{i=1}^n \log(x_i)$$

degree  $\theta = n$

- second-order cone:  $K = \{(x, t) \in \mathbb{R}^n \times \mathbb{R} : \|x\| \leq t\}$

$$\psi(x, t) = \log(t^2 - x^\top x)$$

degree  $\theta = 2$

- positive semidefinite cone:  $K = \mathbb{S}_+^n$

$$\psi(X) = \log \det(X)$$

degree  $\theta = n$

- Consider the approximation

$$\begin{array}{ll} \text{minimize} & tf(x) + \sum_{i=1}^m -\psi_i(-g_i(x)) \\ \text{subject to} & Ax = b \end{array}$$

- Assume an unique minimizer  $x^*(t)$
- $x^*(t)$  is (at most)  $(\sum_{i=1}^m \theta_i)/t$  suboptimal:

$$f(x^*(t)) - p^* \leq \frac{\sum_{i=1}^m \theta_i}{t}$$

- Central path =  $\{x^*(t) : t > 0\}$
- Following the central path leads to a solution

- Example:

$$\begin{array}{ll}\text{minimize} & \text{tr}(AX) \\ \text{subject to} & \text{diag}(X) = 1 \\ & X \succeq 0\end{array}$$

- Log barrier subproblem:

$$\begin{array}{ll}\text{minimize} & t \text{tr}(AX) - \log \det(X) \\ \text{subject to} & \text{diag}(X) = 1\end{array}$$

with implicit constraint  $X \succ 0$

- KKT conditions for subproblem:

$$\begin{cases} tA - X^{-1} + \text{Diag}(\lambda) = 0 \\ \text{diag}(X) = 1 \end{cases}$$

with  $X \succ 0$

- Linearizing around a feasible  $X$ :

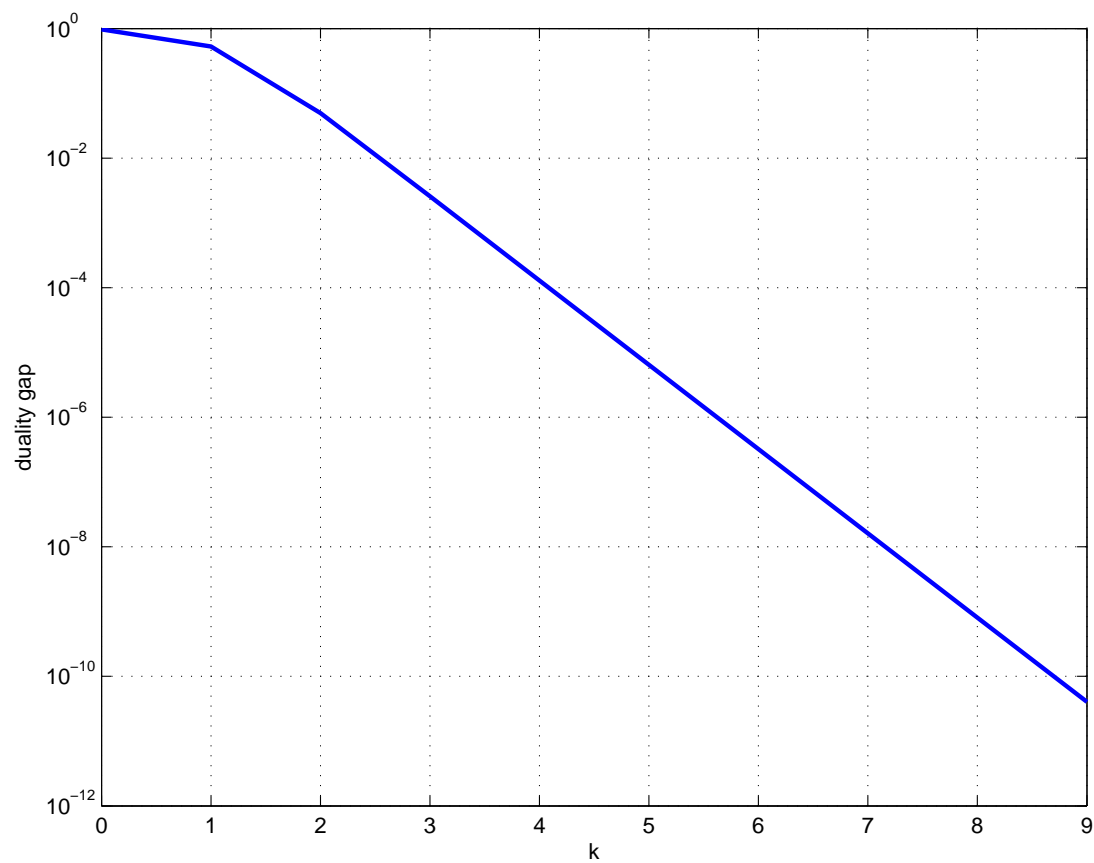
$$\begin{cases} tA - X^{-1} + X^{-1}\Delta X^{-1} + \text{Diag}(\lambda) = 0 \\ \text{diag}(\Delta) = 0 \end{cases}$$

- Solution:

$$\begin{cases} (X \odot X) \lambda = X - tXAX \\ \Delta = X - tXAX - X\text{Diag}(\lambda)X \end{cases}$$

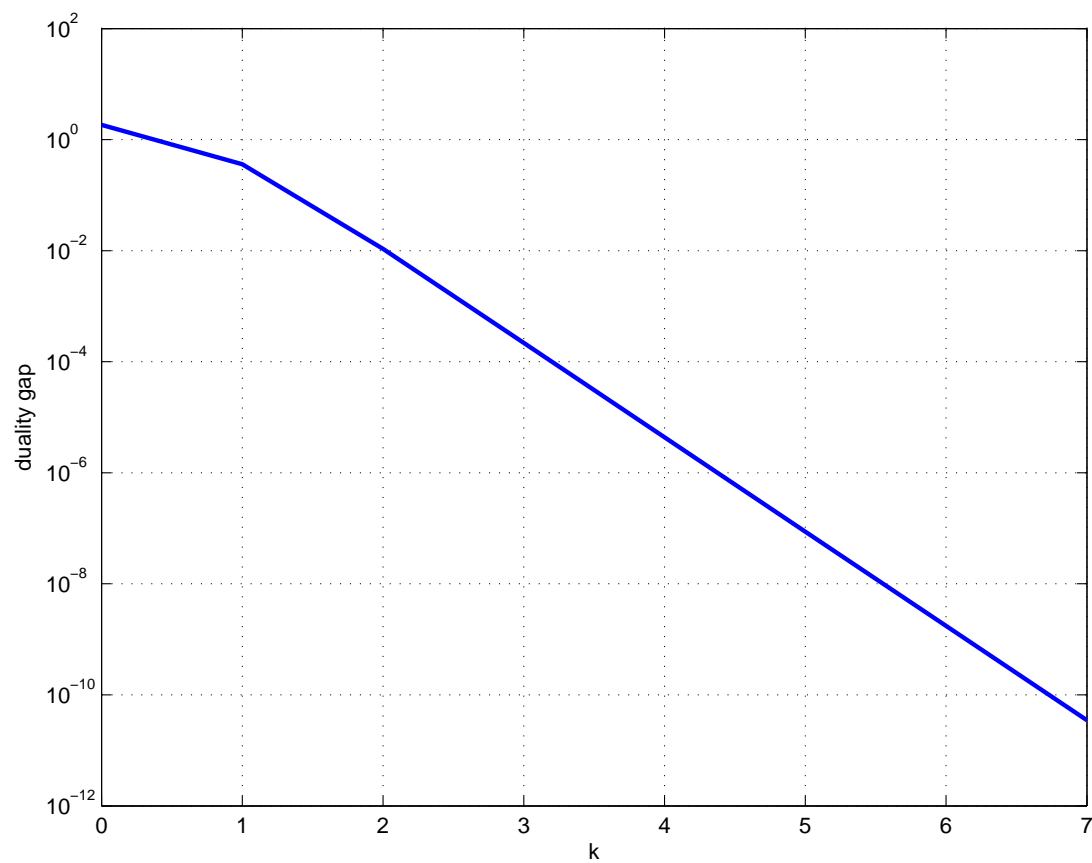


Problem size  $n = 20$  ( $\mu = 20$ )



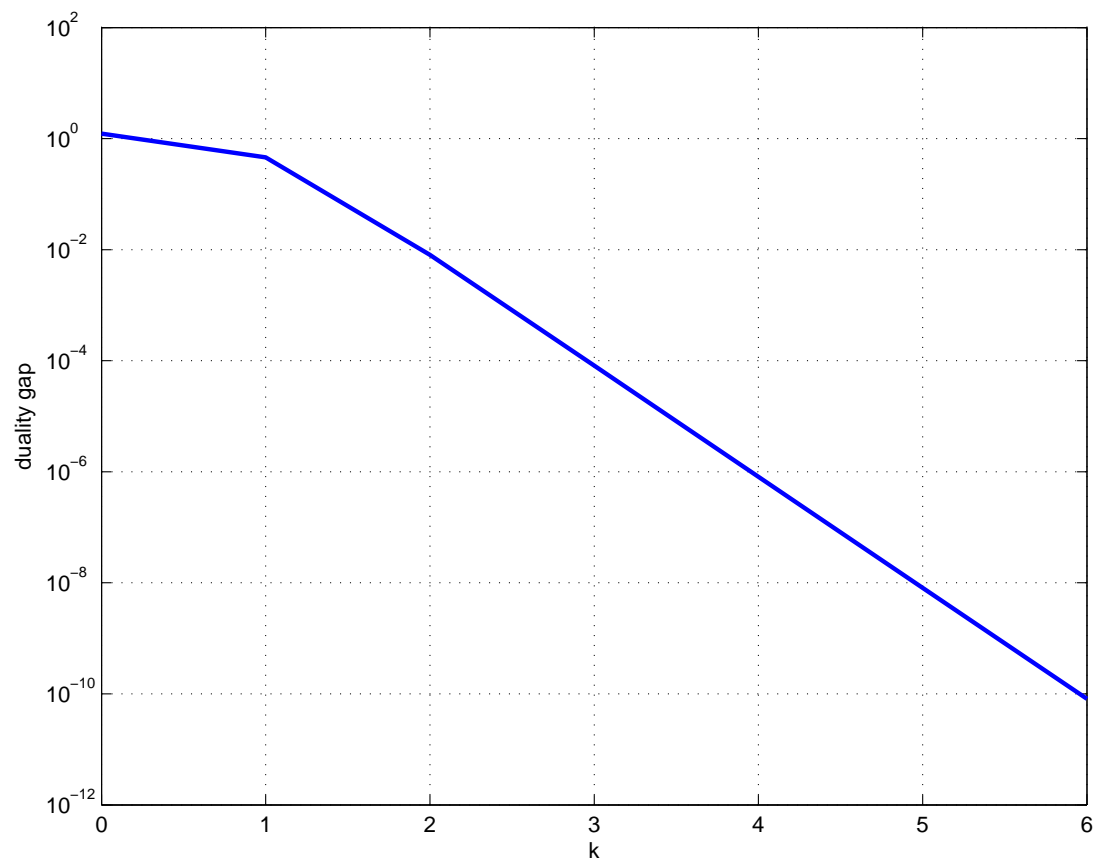
Time  $\simeq 0.06$  sec

Problem size  $n = 20$  ( $\mu = 50$ )



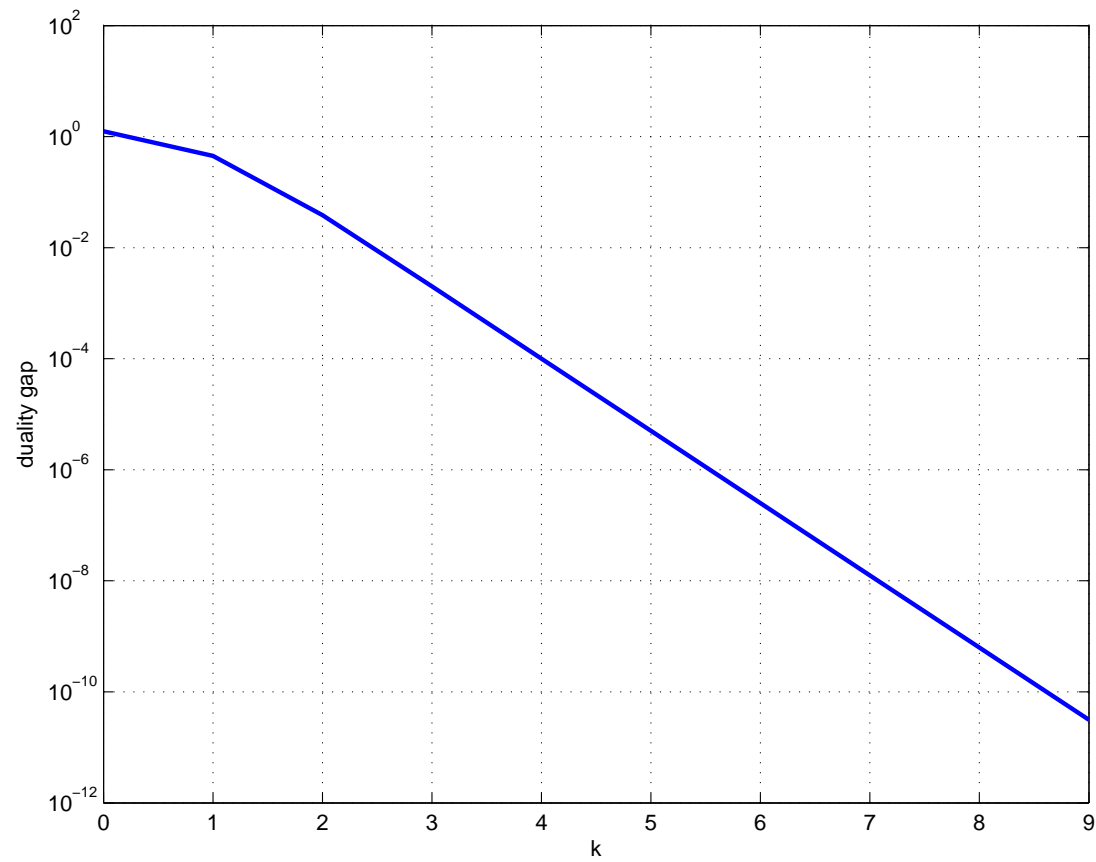
Time  $\simeq 0.06$  sec

Problem size  $n = 20$  ( $\mu = 100$ )



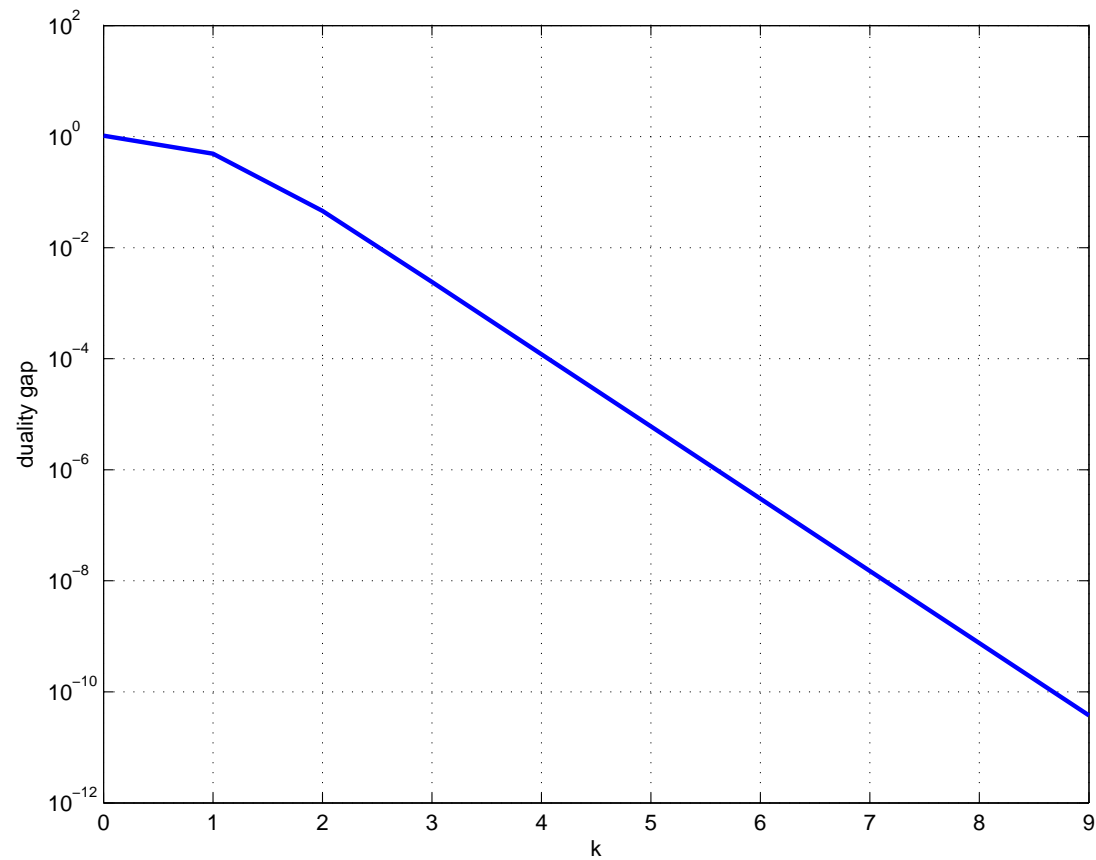
Time  $\simeq 0.06$  sec

Problem size  $n = 100$  ( $\mu = 20$ )



Time  $\simeq 1.2$  sec

Problem size  $n = 150$  ( $\mu = 20$ )



Time  $\simeq 3.4$  sec

Thank you !

Special thanks to **João Paulo** and **Rob**