Modeling and Automation of Industrial Processes

Modelação e Automação de Processos Industriais / MAPI

Supervised Control of Discrete Event Systems Supervision Controllers (Part 2/2)

http://www.isr.tecnico.ulisboa.pt/~jag/courses/mapi2122

Prof. Paulo Jorge Oliveira, original slides Prof. José Gaspar, rev. 2021/2022

Some pointers on Supervised Control of DES

Analysers & simulators

http://www.nd.edu/~isis/techreports/isis-2002-003.pdf (Users Manual)
http://www.nd.edu/~isis/techreports/spnbox/ (Software)

Bibliography:

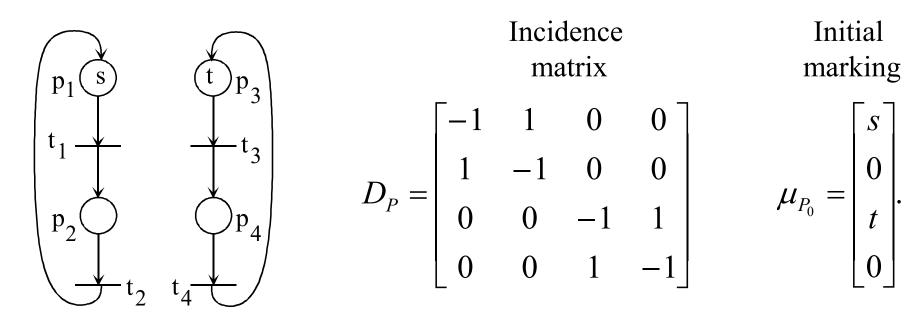
Supervisory Control of Discrete Event Systems using Petri Nets, J. Moody J. and P. Antsaklis, Kluwer Academic Publishers, 1998.

Supervised Control of Concurrent Systems: A Petri Net Structural Approach, M. Iordache and P. Antsaklis, Birkhauser 2006.

Discrete Event Systems - Modeling and Performance Analysis, Christos G. Cassandras, Aksen Associates, 1993.

Feedback Control of Petri Nets Based on Place Invariants, K. Yamalidou, J. Moody, M. Lemmon and P. Antsaklis, http://www.nd.edu/~lemmon/isis-94-002.pdf

Example of controller synthesis: s Producers / t Consumers



Let p2= #machines working, t2= product produced

p3= #consumers, t3= request to consume (e.g. transport product)

Q: How to write *consume only when produced*? What is the linear constraint?

Not possible to write it as a linear constraint on places $L\mu_p \le b$. Is it impossible to solve this problem with the supervised control?

Methods of Synthesis Generalized linear constraint

Let the generalized linear constraint be

$$L\mu_P + Fq_P + Cv_P \leq b,$$

$$\mu_P \in N_0^n, v_P \in N_0^m, q_P \in N_0^m,$$

$$L \in Z^{n_C \times n}, F \in Z^{n_C \times m}, C \in Z^{n_C \times m}, e \quad b \in Z^{n_C},$$

where

- * μ_P is the marking vector for system P
- * q_P is the firing vector since t_0
- * v_P is the number of transitions (firing) that can occur, also designated as Parikh vector

Example: detail elements forming the generalized linear constraint

$$L\mu_P + Fq_P + Cv_P \le b,$$

$$\mu_P \in N_0^n, v_P \in N_0^m, q_P \in N_0^m,$$

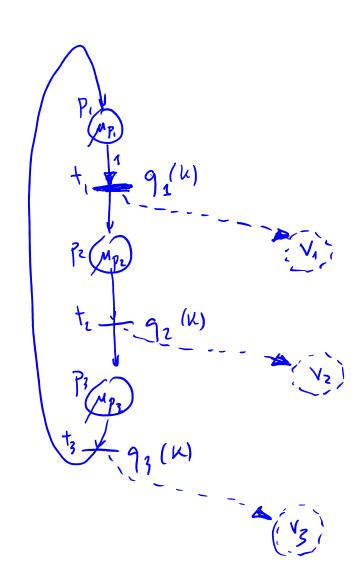
$$L \in Z^{n_C \times n}, F \in Z^{n_C \times m}, C \in Z^{n_C \times m},$$

$$b \in Z^{n_C},$$

State:
$$\mu_p(k) = \begin{bmatrix} \mu_{p1} \\ \mu_{p2} \\ \mu_{p3} \end{bmatrix}_k$$

Firing vector:
$$q_p(k) = \begin{bmatrix} q_1 \\ q_2 \\ q_3 \end{bmatrix}_k$$

Parikh vector:
$$v_p(k) = \begin{bmatrix} v_1 \\ v_2 \\ v_3 \end{bmatrix}_k$$



Function LINENF of SPNBOX

Theorem*: Synthesis of Controllers based on Place Invariants, for **Generalized Linear Constraints**

Given the generalized linear constraint $L\mu_P + Fq_P + Cv_P \le b$, if $b - L\mu_{P_0} \ge 0$, then the controller with incidence matrix and initial marking, respectively

$$\begin{split} D_C^- &= \max\left(0, LD_P + C, F\right) \\ D_C^+ &= \max\left(0, F - \max\left(0, LD_P + C\right)\right) - \min\left(0, LD_P + C\right), \end{split}$$

$$\mu_{C_0} = b - L\mu_{P_0} - Cv_{P_0},$$

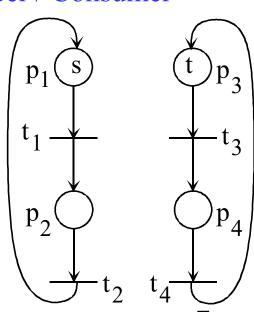
guarantees that constraints are verified for the states resulting from the initial marking.

^{*} In the next slides this will be called the LINENF theorem.

Example of controller synthesis

$$\forall s \in N_0, \forall t \in N_0, \forall n \in N_0$$





Linear constraint: $v_3 \le v_2$

that can be written as:

$$\begin{vmatrix} Cv_P \le b \\ L = 0, F = 0 \end{vmatrix} \begin{bmatrix} 0 & -1 & 1 & 0 \end{bmatrix} \begin{vmatrix} v_1 \\ v_2 \\ v_3 \\ v_4 \end{vmatrix} \le 0.$$

$$D_{P} = \begin{bmatrix} -1 & 1 & 0 & 0 \\ 1 & -1 & 0 & 0 \\ 0 & 0 & -1 & 1 \\ 0 & 0 & 1 & -1 \end{bmatrix} \qquad \begin{array}{c} \text{Initial} \\ \text{marking} \\ \end{array}$$

$$\mu_{P_0} = egin{bmatrix} s \ 0 \ t \ 0 \end{bmatrix}.$$

Example of controller synthesis

Producer / Consumer

$$b - L\mu_{P_0} = 0 - 0 \ge 0.$$

OK.

$$\begin{aligned} D_C^- &= \max \left(0, LD_P + C, F \right) \\ D_C^+ &= \max \left(0, F - \max \left(0, LD_P + C \right) \right) - \min \left(0, LD_P + C \right), \end{aligned}$$

$$D_C^- = \max(0, [0 -1 \ 1 \ 0], \ 0) = [0 \ 0 \ 1 \ 0]$$

$$D_C^+ = \max(0, -[0 \ 0 \ 1 \ 0]) - \min(0, [0 \ -1 \ 1 \ 0])$$

$$= [0 \ 0 \ 0 \ 0] - [0 \ -1 \ 0 \ 0] = [0 \ 1 \ 0 \ 0]$$

$$\mu_{C_0} = b - L\mu_{P_0} - Cv_{P_0},$$

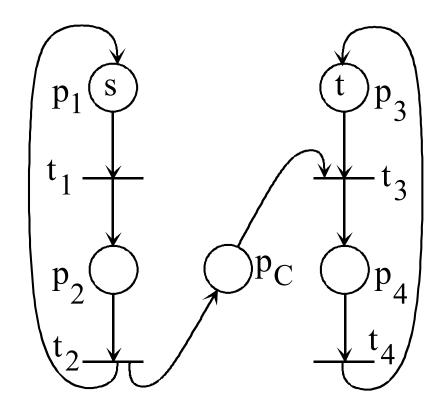
$$\mu_{C_0} = b - L\mu_{P_0} = 0 - 0 = 0.$$

OK.

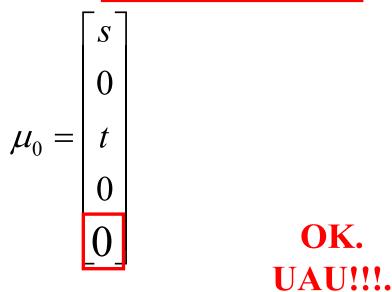
Example of controller synthesis

Producer / Consumer

3) Resulting in



$$D = \begin{bmatrix} -1 & 1 & 0 & 0 \\ 1 & -1 & 0 & 0 \\ 0 & 0 & -1 & 1 \\ 0 & 0 & 1 & -1 \\ \hline 0 & 1 & -1 & 0 \end{bmatrix}$$



```
% The Petri net D=Dp-Dm, and mO
% (Dplus-Dminus= Post-Pre)
Dm= [1 0 0 0;
     0 1 0 0;
     0 0 1 0;
     0 0 0 1];
 Dp= [0 1 0 0;
     1 0 0 0;
     0 0 0 1;
     0 0 1 0];
m0= [1 0 1 0]';
% Supervisor constraint
20
L= []; F= []; C= [0 -1 1 0];
b = 0:
% Computing the supervisor
20
[Dfm, Dfp, ms0] = linenf(Dm, Dp, L, b, m0, F, C)
Df= Dfp-Dfm
ms0
```

Example of controller synthesis: Producer Consumer

Result using the function LINENF.m of the toolbox SPNBOX:

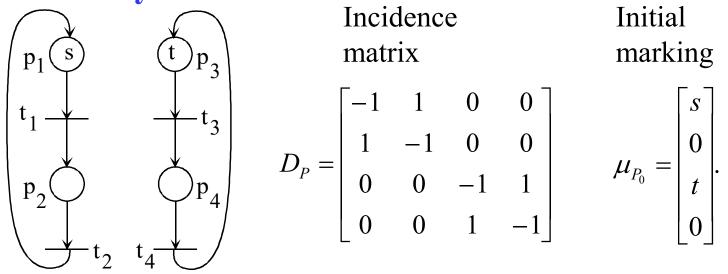
```
Df =
```

```
ms0 =

1
0
1
0
0
```

Example of controller synthesis

Bounded Producer / Consumer



Incidence

$$D_P = \begin{bmatrix} -1 & 1 & 0 & 0 \\ 1 & -1 & 0 & 0 \\ 0 & 0 & -1 & 1 \\ 0 & 0 & 1 & -1 \end{bmatrix}$$

Initial

$$\mu_{P_0} = egin{bmatrix} s \ 0 \ t \ 0 \end{bmatrix}.$$

TWO linear constraints:

$$\begin{cases} v_3 \le v_2 \\ v_2 \le v_3 + n \end{cases} \Leftrightarrow \begin{cases} v_3 - v_2 \le 0 \\ v_2 - v_3 \le n \end{cases}$$

$$\forall s \in N_0, \forall t \in N_0, \forall n \in N_0$$

The two linear constraints can be written as:

$$\begin{cases} v_3 \leq v_2 \\ v_2 \leq v_3 + n \end{cases} \Leftrightarrow \begin{cases} v_3 - v_2 \leq 0 \\ v_2 - v_3 \leq n \end{cases} \qquad Cv_p \leq b \\ i.e. \ L = 0, F = 0 \end{cases} \Leftrightarrow \begin{bmatrix} 0 & -1 & 1 & 0 \\ 0 & 1 & -1 & 0 \end{bmatrix} \begin{bmatrix} v_1 \\ v_2 \\ v_3 \\ v_4 \end{bmatrix} \leq \begin{bmatrix} 0 \\ n \end{bmatrix}$$

Example of controller synthesis

Bounded Producer / Consumer

1) Test

$$b-L\mu_{P_0}=b=\begin{bmatrix}0\\n\end{bmatrix}\geq 0.$$
 OK.

2) Compute

$$\begin{split} D_C^- &= \max \Biggl(0, \begin{bmatrix} 0 & -1 & 1 & 0 \\ 0 & 1 & -1 & 0 \end{bmatrix}, 0 \Biggr) = \begin{bmatrix} 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix}, \\ D_C^+ &= \max \Biggl(0, 0 - \max \Biggl(0, \begin{bmatrix} 0 & -1 & 1 & 0 \\ 0 & 1 & -1 & 0 \end{bmatrix} \Biggr) \Biggr) - \min \Biggl(0, \begin{bmatrix} 0 & -1 & 1 & 0 \\ 0 & 1 & -1 & 0 \end{bmatrix} \Biggr) \\ &= \begin{bmatrix} 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix} - \begin{bmatrix} 0 & -1 & 0 & 0 \\ 0 & 0 & -1 & 0 \end{bmatrix} = \begin{bmatrix} 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \end{bmatrix}. \end{split}$$

and

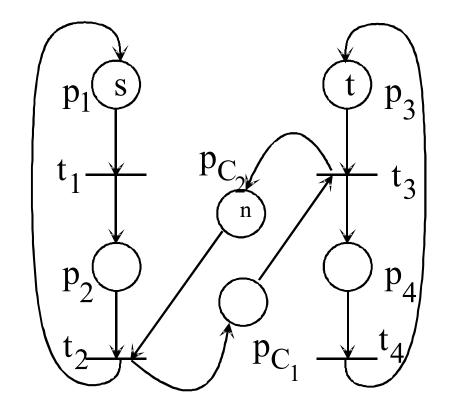
$$\mu_{C_0} = b - L\mu_{P_0} = \begin{vmatrix} 0 \\ n \end{vmatrix}.$$

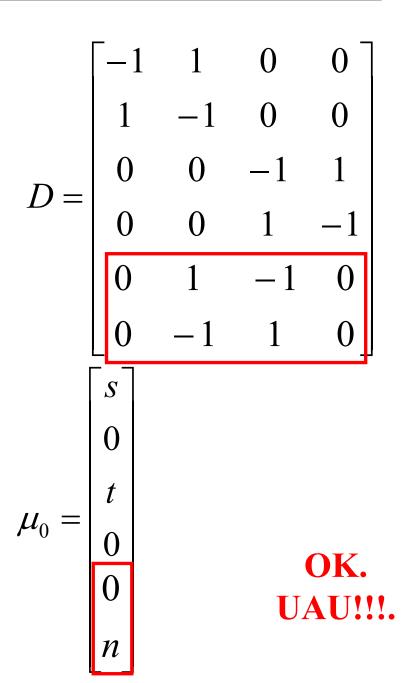
OK.

Example of controller synthesis

Bounded Producer / Consumer

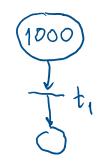
3) Resulting in





Example of controller synthesis – Flow regulation

Consider a Petri net with a large initial marking

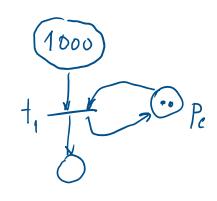


Objective: do *NOT allow consuming* too many tokens in a single step.

For example, one wants to enforce max q_1 to be 2, i.e. accepting only $q_1 = 0$ or $q_1 = 1$ or $q_1 = 2$.

Constraint Solution

1.9,
$$\leq 2$$
 $D_c^{\dagger} = 1$
 $D_c = 1$
 $D_c = 1$
 $D_c = 1$



Function LINENF of SPNBOX

LINENF Lemma 1: From General Constraints to Theorem T1

Given the generalized linear constraint $L\mu_P + Fq_P + Cv_P \le b$ and the conditions of the LINENF theorem:

If
$$L \neq 0$$
, $F = 0$, $C = 0$

then
$$D_C^+ = (LD_P)^-, \quad D_C^- = (LD_P)^+$$
 and $D_C = -LD_P$

$$\mu_{C0} = b - L\mu_{P0}$$

(see proof in the next page)

Notation:

$$D^{+} = \max(0, D)$$

$$D^{-} = -\min(0, D)$$

$$D^{-} = D^{+} - D^{-}$$

$$D^{+}, D^{-} \in N_{0}^{n \times m} \text{ and } D \in Z_{0}^{n \times m}$$

IST / DEEC / API

$$D_{C}^{-} = \max(0, LD_{P} + C, F)$$

$$D_{C}^{+} = \max(0, F - \max(0, LD_{P} + C)) - \min(0, LD_{P} + C),$$

$$\mu_{C_{0}} = b - L\mu_{P_{0}} - Cv_{P_{0}},$$

$$L \neq 0$$
, $F = 0$, $C = 0$ $\Rightarrow L_{MP} \leq b$

$$D_{c} = \max(0, LD_{p} + 1, f)$$

$$= \max(0, LD_{p} + 1, f)$$

$$= \max(0, LD_{p})$$

$$= (LD_{p})^{\dagger}$$

$$= (LD_{p})^{\dagger}$$

$$= + (LD_{p})^{\dagger}$$

$$= + (LD_{p})^{\dagger}$$

$$D_{c} = D_{c}^{+} - D_{c}^{-} = (LD_{p})^{-} - (LD_{p})^{+} = -((LD_{p})^{+} - (LD_{p})^{-}) = -LD_{p}/$$

$$M_{co} = b - LM_{po} - 4V_{po} = b - LM_{po}/$$

Function LINENF of SPNBOX

LINENF Lemma 2: Firing Regulation

Given the generalized linear constraint $L\mu_P + Fq_P + Cv_P \le b$ and the conditions of the LINENF theorem:

If
$$L=0$$
, $F\neq 0$, $C=0$ then $D_C^+=F^+$, $D_C^-=F^+$ and $D_C=0$
$$\mu_{C0}=b$$

(homework, prove this lemma)

Function LINENF of SPNBOX

LINENF Lemma 3: Constraints on Counters

Given the generalized linear constraint $L\mu_P + Fq_P + Cv_P \le b$ and the conditions of the LINENF theorem:

If
$$L = 0$$
, $F = 0$, $C \neq 0$

then
$$D_C^+ = C^-, \qquad D_C^- = C^+$$

$$D_C^- = C^+$$

and

$$D_C = -C$$

$$\mu_{C0} = b - C \nu_{P0}$$

(homework, prove this lemma)

(empty page, do yourself the proof of the last two lemmas)

Methods of Synthesis: intro to Uncontrollable and Unobservable transitions

Definition of Uncontrollable Transition:

A transition is uncontrollable if its firing **cannot be inhibited** by an external action (e.g. a supervisory controller).



Definition of Unobservable Transition:

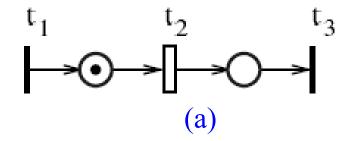
A transition is unobservable if its firing cannot be detected or measured (therefore the study of any supervisory controller can not depend from that firing).

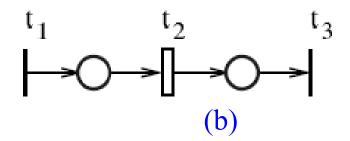


Proposition:

In a Petri net based controller, both input and output arcs to/from plant transitions are used to trigger state changes in the controller. Since a controller cannot have arcs connecting to unobservable transitions, then all unobservable transitions are also implicitly uncontrollable.

Methods of Synthesis: intro to Uncontrollable and Unobservable transitions





If t1 is controllable and t2 is uncontrollable:

- case (a), then t2 cannot be directly inhibited; it will eventually fire
- case (b), then t2 can be indirectly prevented from firing by inhibiting t1.

i.e. may exist indirect solution despite t2 being uncontrollable.

If <u>t2 is unobservable</u> and <u>t3 is observable</u>, then we cannot detect when t2 fires. The state of a supervisor is not changed by firing t2. However we can <u>indirectly detect that t2 has fired</u>, by detecting the firing of t3.

i.e. may exist indirect solution despite t2 being unobservable.

∴ may exist indirect solution despite t2 uncontrollable and/or unobservable.

Definition: A marking μ_P is admissible if

i) $L\mu_P \le b$ and ii) $\forall \mu' \in R(C, \mu_P)$ verifies $L\mu' \le b$

Definition: A Linear Constraint (L, b) is admissible if

- i) $L\mu_{Po} \leq b$ and
- ii) $\forall \mu' \in R(C, \mu_{Po})$ such that $L\mu' \leq b$ μ' is an admissible marking.

Note: ii) indicates that the firing of uncontrollable transitions can never lead from a state that satisfies the constraint to a new state that does not satisfy the constraint.

Proposition: Admissibility of a constraint

A linear constraint is admissible *iff*

- The initial markings satisfy the constraint.
- There exists a controller with maximal permissivity that forces the constraint and does not inhibit any uncontrollable transition.

Two sufficient (not necessary) conditions:

Corollary: given a system with uncontrollable transitions, $l^T D_{uc} \leq 0$ implies admissibility.

Corollary: given a system with unobservable transitions, $l^T D_{uo} = 0$ implies admissibility.

Function MRO ADM of SPNBOX

Lemma *: Structure of Constraint transformation

If
$$L'\mu_p \leq b'$$
 where

$$L' = R_1 + R_2L$$
 and $b' = R_2(b+1) - 1$

$$R_1 \in Z^{n_C \times n}$$
 and $R_1 \mu_p \ge 0$

 $R_2 \in \mathbb{Z}^{n_C \times n_C}$ is a matrix with positive elements in the diagonal

then

$$L\mu_p \le b$$

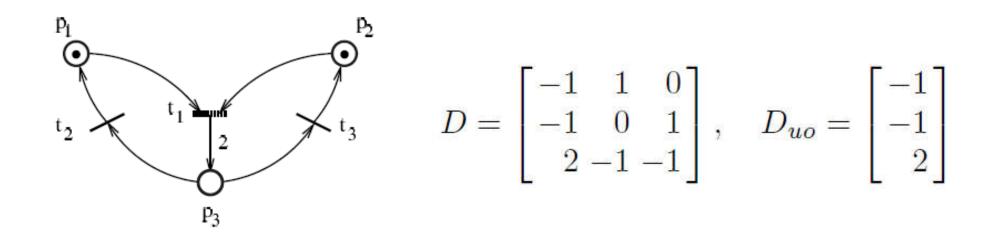
 $L\mu_p \le b$ is also verified.

Typical usage:

- constraint (L, b), and extra constraints $\rightarrow (R_1, R_2) \rightarrow (L', b') \rightarrow (D_C^+, D_C^-, \mu_{C0})$
- see unobservable / uncontrollable extra constraints example in the next slides

^{*} Lemma 4.10 in [Moody98] pg46

Example: design controller with t1 unobservable (1/4)



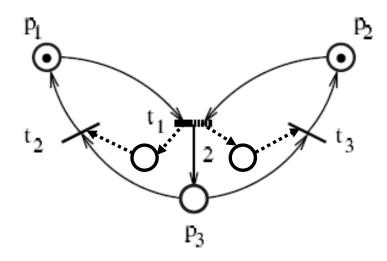
Objectives: $\mu_1 + \mu_3 \ge 1$ and $\mu_2 + \mu_3 \ge 1$ which can be written in matrix form as

$$L\mu \le b$$
, $L = \begin{bmatrix} -1 & 0 & -1 \\ 0 & -1 & -1 \end{bmatrix}$, $b = \begin{bmatrix} -1 \\ -1 \end{bmatrix}$

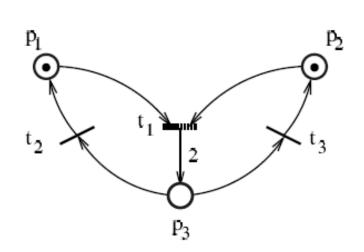
Example: design controller with t1 unobservable (2/4)

```
% System and constraints
D = [-1 \ 1 \ 0;
    -1 0 1;
    +2 -1 -1];
Dm = -D.*(D<0);
Dp = D.*(D>0);
m0= [1 1 0]';
L= [-1 \ 0 \ -1; \ 0 \ -1 \ -1];
b = [-1; -1];
% Supervisor computation
[Dfp, Dfm, mf0] =
    linenf( Dp, Dm, L, b, m0 );
```

 $^{\land}$ Bad news, supervisor touches t_1 .



Example: design controller with t1 unobservable (3/4)



[La, ba, R1, R2] = mro_adm(L, b, D, Tuc, Tuo);

Solution obtained with the function MRO ADM.m of the SPNBOX toolbox:

$$R1 = \begin{bmatrix} 0 & 0 & 1 \\ 0 & 0 & 1 \end{bmatrix}$$

$$\mathbf{La} = \begin{bmatrix} -2 & 0 & -1 \\ 0 & -2 & -1 \end{bmatrix} \qquad \mathbf{ba} = \begin{bmatrix} -1 \\ -1 \end{bmatrix}$$

$$ba = \begin{bmatrix} -1 \\ -1 \end{bmatrix}$$

$$R2 = \begin{bmatrix} 2 & 0 \\ 0 & 2 \end{bmatrix}$$

Note: verify that $L_a \mu \leq b_a$ implies $L \mu \leq b$

Example: design controller with t1 unobservable (4/4)

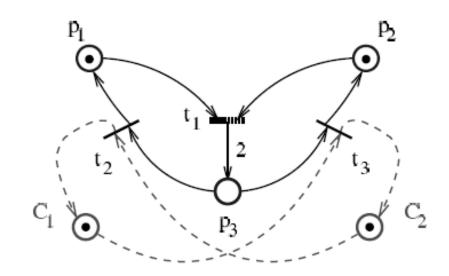
Finally the supervised controller is simply obtained from L_a and b_a :

$$\begin{split} D_c &= -L_a D_p \\ &= \begin{bmatrix} -2 & 0 & -1 \\ 0 & -2 & -1 \end{bmatrix} \begin{bmatrix} -1 & 1 & 0 \\ -1 & 0 & 1 \\ 2 & -1 & -1 \end{bmatrix} \\ &= \begin{bmatrix} 0 & 1 & -1 \\ 0 & -1 & 1 \end{bmatrix} \end{split}$$

$$\mu_{c0} = b_a - L_a \mu_{p0}$$

$$= \begin{bmatrix} -1 \\ -1 \end{bmatrix} - \begin{bmatrix} -2 & 0 & -1 \\ 0 & -2 & -1 \end{bmatrix} \begin{bmatrix} 1 \\ 1 \\ 0 \end{bmatrix}$$

$$= \begin{bmatrix} 1 \\ 1 \end{bmatrix}$$



Obtained the desired result: supervisor does not touch t_1 .

End of chapter on supervision control.

What is next?