Industrial Automation (Automação de Processos Industriais)

Supervised Control of Discrete Event Systems

http://users.isr.ist.utl.pt/~jag/courses/api1718/api1718.html

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Syllabus:

•••

Chap. 8 - DESs and Industrial Automation [2 weeks]

Chap. 9 – Supervised Control of DESs [1 week] * SCADA * Methodologies for the Synthesis of Supervision Controllers * Failure detection

Some jokes available in <u>http://members.iinet.net.au/~ianw/cartoon.html</u> The End.

Some pointers on Supervised Control of DES

History:	The SCADA Web, http://members.iinet.net.au/~ianw/
	Monitoring and Control of Discrete Event Systems, Stéphane Lafortune,
	http://www.ece.northwestern.edu/~ahaddad/ifac96/introductory_workshops.html
Tutorial:	http://vita.bu.edu/cgc/MIDEDS/ http://www.daimi.au.dk/PetriNets/
Analysers &	http://www.nd.edu/~isis/techreports/isis-2002-003.pdf (Users Manual)
Simulators:	http://www.nd.edu/~isis/techreports/spnbox/ (Software)
Bibliography:	* SCADA books <u>http://www.sss-mag.com/scada.html</u>
815	* K. Stouffer, J. Falco, K. Kent, "Guide to Supervisory Control and Data
	Acquisition (SCADA) and Industrial Control Systems Security",
	NIST Special Publication 800-82, 2006
	* Moody J. e Antsaklis P., "Supervisory Control of Discrete Event
	Systems using Petri Nets," Kluwer Academic Publishers, 1998.
	* Cassandras, Christos G., "Discrete Event Systems - Modeling and
	Performance Analysis," Aksen Associates, 1993.
	* Yamalidou K., Moody J., Lemmon M. and Antsaklis P.
	Feedback Control of Petri Nets Based on Place Invariants
	http://www.nd.edu/~lemmon/isis-94-002.pdf

Chap. 9 – Supervised Control of DES

Supervision of DES: SCADA

Supervisory

Control

And

Data

Acquisition

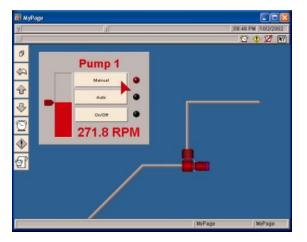
IST / DEEC / API

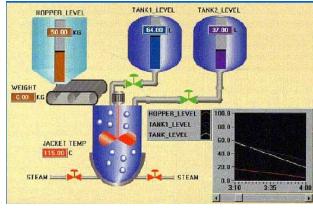
Chap. 9 – Supervised Control of DES

Supervision of DES SCADA interface / GUI



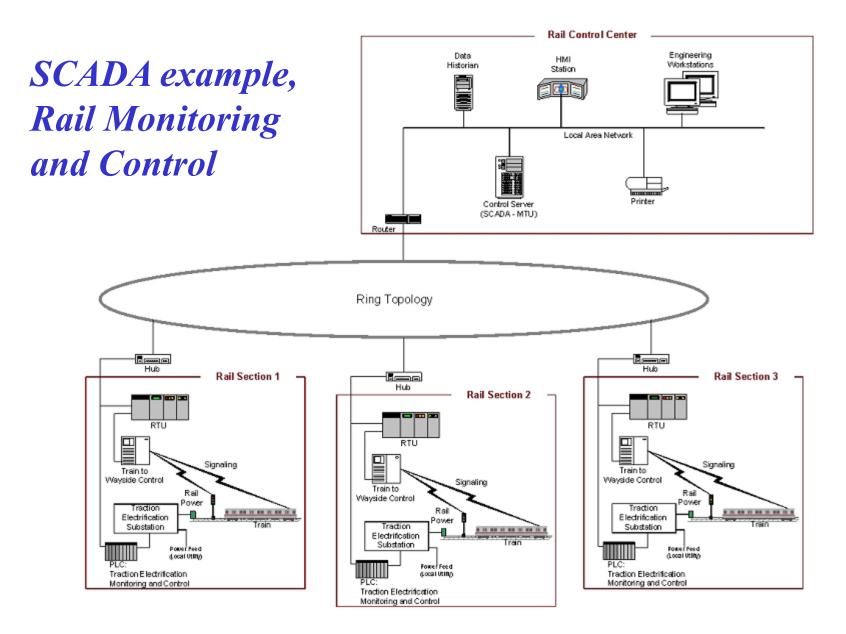








IST / DEEC / API



SCADA vs ICS

Industrial Control Systems (ICS):

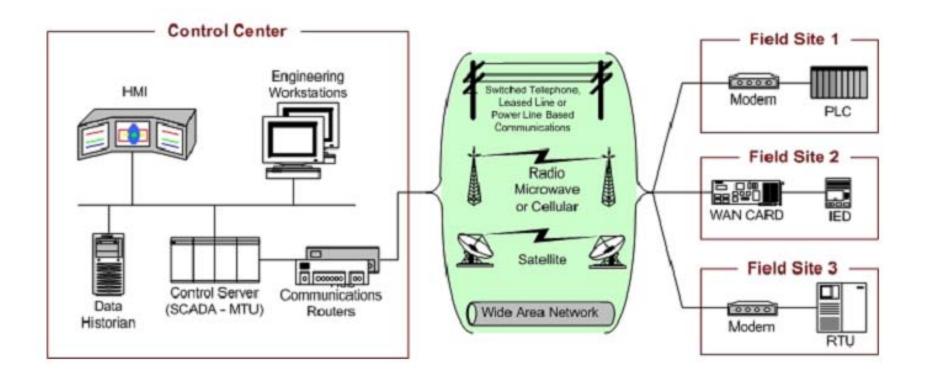
- Supervisory Control and Data Acquisition (SCADA) systems,
- Distributed Control Systems (DCS), or
- smaller configurations such as skid-mounted PLCs

ICSs are typically used in industries such as electric, water, oiland-gas, transportation, chemical, pharmaceutical, pulp-andpaper, food and beverage, and discrete-manufacturing (e.g. automotive, aerospace, and durable goods).

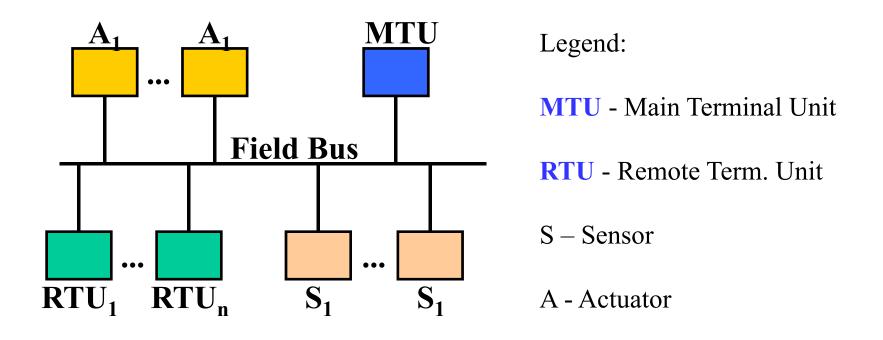
SCADA topics

- Remote monitoring of the state of automation systems
- Logging capacity (resorting to specialized Databases)
- Able to access to *historical* information (plots along time, with selectable periodicity)
- Advanced tools to design Human-Machine interfaces
- Failure Detection and Isolation capacity (*threshold* and/or logical functions) on supervised quantities
- Access control

SCADA system general layout



Hardware Support Architecture of SCADA



More terminology: Field Bus (IEC 61158) / PROFIBUS (Field Bus type, Siemens), MODBUS (Schneider)

Examples of software packages including SCADA solutions

- Aimax, de Desin Instruments S.A.
- CUBE, Orsi España S.A.
- FIX, de Intellution.
- Lookout, National Instruments.
- Monitor Pro, de Schneider Electric.
- SCADA InTouch, de LOGITEK.
- SYSMAC SCS, de Omron.
- Scatt Graph 5000, de ABB.
- WinCC, de Siemens.

IST / DEEC / API

Chap. 9 – Supervised Control of DES

Supervision of DES

And

Now

Something

Completely

Different

Objectives of the Supervised Control

- Supervise and bound the work of the supervised DES
- Reinforce that some properties are verified
- Assure that some states are not reached
- Performance criteria are verified
- Prevent deadlocks in DES
- Constrain on the use of resources (e.g. mutual exclusion)

Some history on Supervised Control

- Methods for finite automata [Ramadge et *al.*], 1989
 - some are based on brute-force search (!)
 - or may require simulation (!)

• Formal verification of *software* in Computer Science (since the 60s) and on *hardware* (90, ...)

• Supervisory Control Method of Petri Nets, method based on *monitors* [Giua et *al*.], 1992.

• Supervisory Control of Petri Nets based on Place Invariants [Moody, Antsaklis et *al*.], 1994 (shares some similarities with the previous one, but deduced independently!...).

Advantages of the Supervisory Control of Petri Nets

- Mathematical representation is clear (and easy)
- Resorts only to linear algebra (matrices)
- More compact then automata
- Straightforward the representation of infinity state spaces
- Intuitive graphical representation available

The representation of the controller as a Petri Net leads to simplified Analysis and Synthesis tasks

Place Invariants

Place invariants are sets of places whose token count remains always constant. Place invariants can be computed from integer solutions of $w^T D = 0$. Non-zero entries of w correspond to the places that belong to the particular invariant.

Supervisor Synthesis using Place Invariants [ISIS docs]:

What type of relations can be represented in the method of Place Invariants?

- Sets of linear constraints in the state space
- Representation of convex regions (there are extensions for non-convex regions)
- Constraints to guarantee liveness and to avoid deadlocks (that can be expressed, in general, as linear constraints)
- Constraints on the events and timings (that can be expressed, in general, as linear constraints)

Methods of Analysis/Synthesis

Method of the Matrix Equations (just to remind)

The dynamics of the Petri net state can be written in compact form as:

$$\mu(k+1) = \mu(k) + Dq(k)$$

where:

- $\mu(k+1)$ marking to be reached
- μ (k) initial marking
- q(k) firing vector (transitions)
- D incidence matrix. Accounts the balance of tokens, giving the transitions fired.

Methods of Analysis/Synthesis

How to build the Incidence Matrix? (just to remind)

For a Petri net with *n* places and *m* transitions

$$\mu \in N_0^{n}$$

$$q \in N_0^{m}$$

$$D = D^+ - D^-, \quad D \in Z^{n \times m}, \quad D^+ \in N_0^{n \times m}, \quad D^- \in N_0^{n \times m}$$
The enabling firing rule is
$$\mu \ge D^- q$$

Can also be written in compact form as the inequality $\mu + Dq \ge 0$, interpreted element-by-element.

Note: in this course all vector and matrix inequalities are read element-by-element unless otherwise stated.

Some notation for the method

• The supervised system is modelled as a Petri net with *n* places and *m* transitions, and incidence matrix

$$D_P \in \mathbb{Z}^{n \times m}.$$

• The supervisor is modelled as a Petri net with n_C places and m transitions, and incidence matrix

$$D_C \in \mathbb{Z}^{n_C \times m}.$$

• The resulting total system has an incidence matrix

$$D \in \mathbb{Z}^{(n+n_C) \times m}.$$

Theorem: Synthesis of Controllers based on Place Invariants (T1)

Given the set of linear state constraints that the supervised system must follow, written as

$$L\mu_P \leq b, \quad \mu_P \in N_0^n, \quad L \in Z^{n_C \times n} \quad and \quad b \in Z^{n_C}.$$

If
$$b-L\mu_{P_0}\geq 0$$
,

then the controller with incidence matrix and the initial marking, respectively

$$D_C = -LD_P$$
, and $\mu_{C_0} = b - L\mu_{P_0}$,

enforce the constraints to be verified for all markings obtained from the initial marking.

Theorem - proof outline :

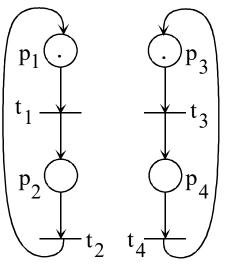
The constraint $L\mu_P \leq b$ can be written as $L\mu_P + \mu_C = b$, using the slack variables μ_C . They represent the marking of the n_C places of the controller.

To have a place invariant, the relation $w^T D = 0$ must be verified and in particular, given the previous constraint:

$$w^T D = \begin{bmatrix} L & I \end{bmatrix} \begin{bmatrix} D_P \\ D_C \end{bmatrix} = 0$$
, resulting $D_C = -LD_P$.

From $L\mu_{P_0} + \mu_{C_0} = b$, follows that $\mu_{C_0} = b - L\mu_{P_0}$.

Example of controller synthesis: Mutual Exclusion



Linear constraint: $\mu_2 + \mu_4 \le 1$ that can be written as: $\begin{aligned} L\mu_P \leq b & \begin{bmatrix} 0 & 1 & 0 & 1 \end{bmatrix} \begin{vmatrix} \mu_1 \\ \mu_2 \\ \mu_3 \\ \mu_4 \end{vmatrix} \leq 1. \end{aligned}$ $D_{P} = \begin{bmatrix} -1 & 1 & 0 & 0 \\ 1 & -1 & 0 & 0 \\ 0 & 0 & -1 & 1 \\ 0 & 0 & 1 & -1 \end{bmatrix} \quad \text{and initial} \quad \mu_{P_{0}} = \begin{bmatrix} 1 \\ 0 \\ 1 \\ 0 \end{bmatrix}.$

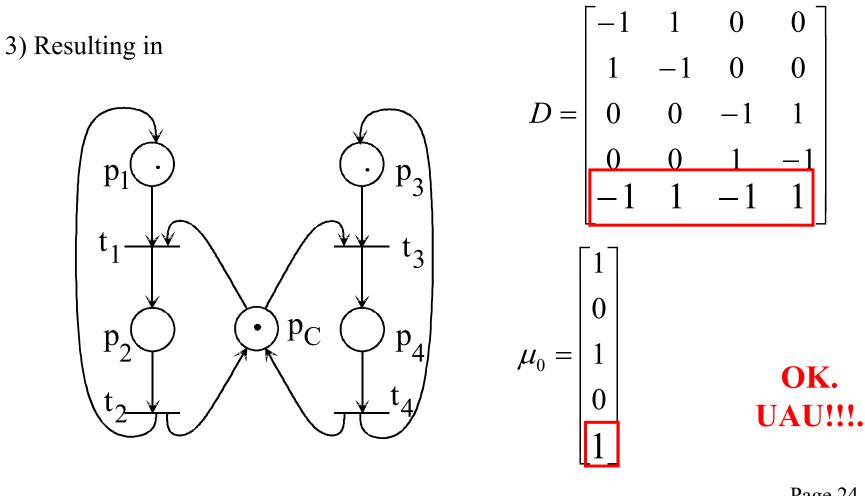
Incidence Matrix

Example of controller synthesis: Mutual Exclusion

1) Test $b - L\mu_{P_0} = 1 - \begin{bmatrix} 0 & 1 & 0 & 1 \end{bmatrix} \begin{bmatrix} 1 & 0 \\ 1 & 0 \\ 1 & 0 \end{bmatrix} = 1 \ge 0.$ 2) Compute $D_C = -LD_P = -\begin{bmatrix} 0 & 1 & 0 & 1 \end{bmatrix} \begin{bmatrix} -1 & 1 & 0 & 0 \\ 1 & -1 & 0 & 0 \\ 0 & 0 & -1 & 1 \\ 0 & 0 & 1 & -1 \end{bmatrix} = \begin{bmatrix} -1 & 1 & -1 & 1 \end{bmatrix},$ and

$$\mu_{C_0} = b - L \mu_{P_0} = 1.$$
 OK.

Example of controller synthesis: Mutual Exclusion



Example of controller synthesis: Mutual Exclusion

```
% The Petri net D=Dp-Dm, and m0
                                                  Result using the function
% (Dplus-Dminus= Post-Pre)
                                                  LINENF.m of the
Dm= [1 0 0 0;
                                                  toolbox SPNBOX:
    0 1 0 0;
    0 0 1 0;
    0 0 0 1];
                                                  Df =
Dp = [0 \ 1 \ 0 \ 0;
                                                      -1 1
                                                                   0
                                                                        0
    1000;
                                                            -1 0
                                                       1
                                                                        0
    0 0 0 1;
                                                       0 0 -1 1
    0 0 1 0];
                                                           0 1
                                                       0
                                                                       -1
                                                      -1 1
                                                                 -1
                                                                        1
m0 = [1 \ 0 \ 1 \ 0]';
% Supervisor constraint
                                                  ms0 =
옿
L = [0 \ 1 \ 0 \ 1];
                                                       1
b = 1;
                                                       0
                                                       1
% Computing the supervisor
                                                       0
읗
                                                       1
[Dfm, Dfp, ms0] = linenf(Dm, Dp, L, b, m0);
Df= Dfp-Dfm
ms0
```

Definition:

Maximal permissivity occurs when (i) all the linear constraints are verified and (ii) all legal markings can be reached.

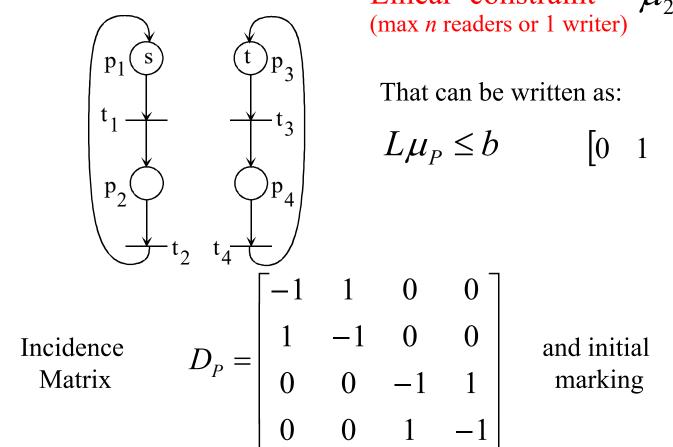
Lemmas:

L1) The controllers obtained with T1 have maximal permissivity.

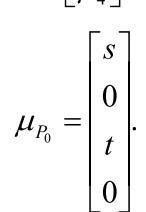
L2) Given the linear constraints used, the place invariants obtained with the controller synthesized with T1 are the same as the invariants associated with the initial system.

Example of controller synthesis

n Readers / 1 Writer



 $\forall s \in N_0, \forall t \in N_0, \forall n \in N_0$ Linear constraint $\mu_2 + n\mu_4 \le n$ (max *n* readers or 1 writer) $L\mu_P \le b$ $\begin{bmatrix} 0 & 1 & 0 & n \end{bmatrix} \begin{bmatrix} \mu_1 \\ \mu_2 \\ \mu_3 \end{bmatrix} \le n.$



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Example of controller synthesis

n Readers / 1 Writer 1) Test $b - L\mu_{P_0} = n - \begin{bmatrix} 0 & 1 & 0 & n \end{bmatrix} \begin{bmatrix} s \\ 0 \\ t \\ 0 \end{bmatrix} = n \ge 0.$ OK.

2) Compute

$$D_{C} = -LD_{P} = -\begin{bmatrix} 0 & 1 & 0 & n \end{bmatrix} \begin{bmatrix} -1 & 1 & 0 & 0 \\ 1 & -1 & 0 & 0 \\ 0 & 0 & -1 & 1 \\ 0 & 0 & 1 & -1 \end{bmatrix} = \begin{bmatrix} -1 & 1 & -n & n \end{bmatrix},$$

and

$$\mu_{C_0} = b - L\mu_{P_0} = n.$$
 OK.

0

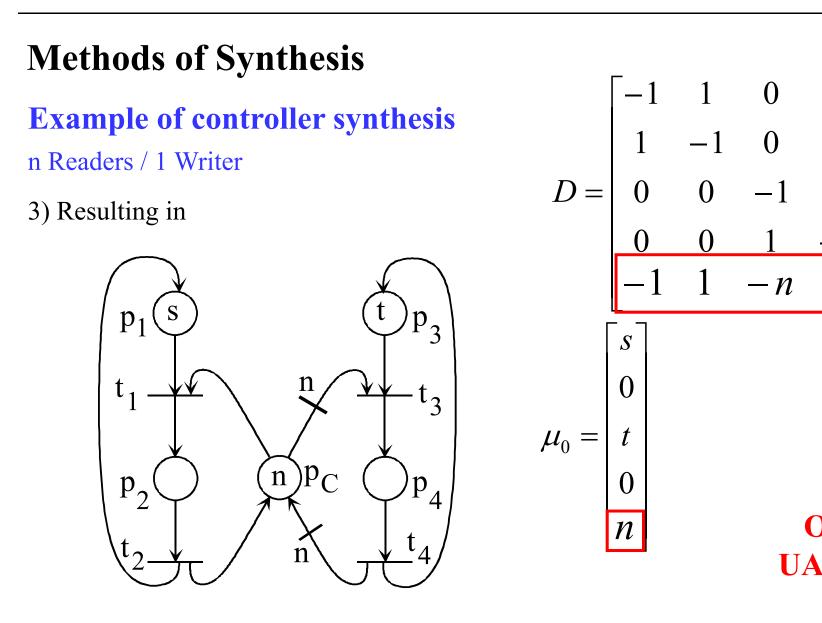
-1

-n

0

0

n





OK.

Advantages of the Method of the Place Invariants [ISIS docs]:

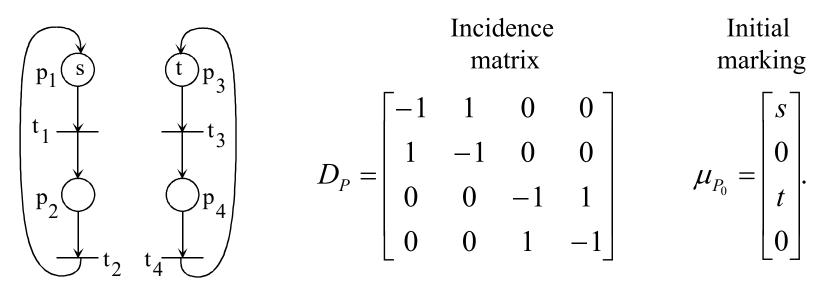
Other characteristics that can impact on the solutions?

- Existence and uniqueness
- Optimality of the solutions (e.g. maximal permissivity)
- Existence of transition non-controllable and/or not observable (remind definitions for time-driven systems)

In general the solutions can be found solving:

Linear Programming Problems, with Linear Constraints

Example of controller synthesis: s **Producers / t Consumers**



 Let p2= #machines working, t2= product produced p3= #consumers, t3= request to consume (e.g. transport product)
 Q: How to write *consume only when produced* ? What is the linear constraint?

Not possible to write it as a linear constraint on places $L\mu_p \le b$. Is it impossible to solve this problem with the proposed method?

Generalized linear constraint

Let the generalized linear constraint be

$$\begin{split} L\mu_P + Fq_P + C\nu_P &\leq b, \\ \mu_P &\in N_0^{\ n}, \nu_P \in N_0^{\ m}, q_P \in N_0^{\ m}, \\ L &\in Z^{n_C \times n}, F \in Z^{n_C \times m}, C \in Z^{n_C \times m}, e \quad b \in Z^{n_C}, \end{split}$$

where

- * μ_P is the marking vector for system P;
- * q_P is the firing vector since t_0 ;
- * v_P is the number of transitions (firing) that can occur, also designated as Parikh vector.

Function LINENF of SPNBOX

Theorem: Synthesis of Controllers based on Place Invariants, for Generalized Linear Constraints

Given the generalized linear constraint $L\mu_P + Fq_P + Cv_P \le b$, if $b - L\mu_{P_0} \ge 0$, then the controller with incidence matrix and initial marking, respectively

$$D_{C}^{-} = \max(0, LD_{P} + C, F)$$

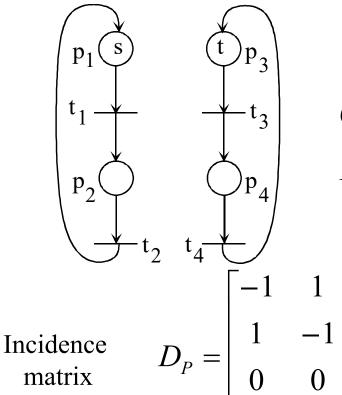
$$D_{C}^{+} = \max(0, F - \max(0, LD_{P} + C)) - \min(0, LD_{P} + C),$$

$$\mu_{C_0} = b - L \mu_{P_0} - C v_{P_0},$$

guarantees that constraints are verified for the states resulting from the initial marking.

Example of controller synthesis $\forall s \in N_0, \forall t \in N_0, \forall n \in N_0$

Producer / Consumer



Linear constraint: $v_3 \leq v_2$ that can be written as: $\begin{vmatrix} \cdot & \cdot \\ - & \cdot \\ D_{P} = \begin{bmatrix} -1 & 1 & 0 & 0 \\ 1 & -1 & 0 & 0 \\ 0 & 0 & -1 & 1 \\ 0 & 0 & 1 & -1 \end{bmatrix}$ Initial marking

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Example of controller synthesis

Producer / Consumer

1) Test
$$b - L\mu_{P_0} = 0 - 0 \ge 0.$$
 OK.

2) Compute

$$D_{C}^{-} = \max(0, LD_{P} + C, F)$$

$$D_{C}^{+} = \max(0, F - \max(0, LD_{P} + C)) - \min(0, LD_{P} + C),$$

$$D_{C}^{-} = \max(0, [0 -1 \ 1 \ 0], \ 0) = [0 \ 0 \ 1 \ 0]$$

$$D_{C}^{+} = \max(0, -[0 \ 0 \ 1 \ 0]) - \min(0, [0 \ -1 \ 1 \ 0])$$

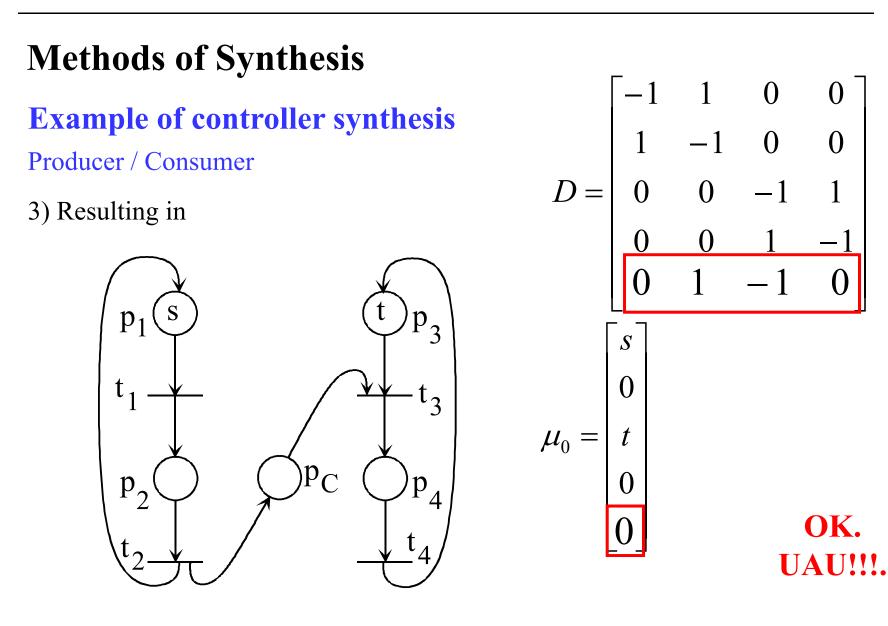
$$= [0 \ 0 \ 0 \ 0] - [0 \ -1 \ 0 \ 0] = [0 \ 1 \ 0 \ 0]$$

and

$$\mu_{C_0} = b - L\mu_{P_0} - Cv_{P_0},$$

$$\mu_{C_0} = b - L\mu_{P_0} = 0 - 0 = 0.$$
 OK.

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% The Petri net D=Dp-Dm, and m0 % (Dplus-Dminus= Post-Pre)

Dm= [1 0 0 0;	
0 1 0 0;	
0 0 1 0;	
0 0 0 1];	
Dp= [0 1 0 0;	
1000;	
0 0 0 1;	
0 0 1 0];	
mO= [1 0 1 0]';	
% Supervisor constraint %	
L= []; F= []; C= [0 -1 1 0];	
b= 0;	
% Computing the supervisor	
*	
[Dfm, Dfp, ms0] = linenf(Dm, Dp, L, b, m0, F,	, C)
Df= Dfp-Dfm	
ns0	

Example of controller synthesis: Producer Consumer

Result using the function LINENF.m of the toolbox SPNBOX:

Df =

ms0 =

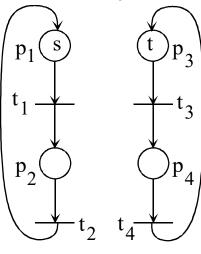
0

-1	1	0	0
1	-1	0	0
0	0	-1	1
0	0	1	-1
0	1	-1	0

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Example of controller synthesis

Bounded Producer / Consumer



 Incidence matrix
 Initial marking

 $D_P = \begin{bmatrix} -1 & 1 & 0 & 0 \\ 1 & -1 & 0 & 0 \\ 0 & 0 & -1 & 1 \\ 0 & 0 & 1 & -1 \end{bmatrix}$ $\mu_{P_0} = \begin{bmatrix} s \\ 0 \\ t \\ 0 \end{bmatrix}$

TWO linear constraints:

$$\begin{cases} v_3 \leq v_2 \\ v_2 \leq v_3 + n \end{cases} \Leftrightarrow \begin{cases} v_3 - v_2 \leq 0 \\ v_2 - v_3 \leq n \end{cases}$$
$$\forall s \in N_0, \forall t \in N_0, \forall n \in N_0 \end{cases}$$

The two linear constraints can be written as:

$$Cv_{P} \leq b$$

i.e. $L = 0, F = 0$ $\Leftrightarrow \begin{bmatrix} 0 & -1 & 1 & 0 \\ 0 & 1 & -1 & 0 \end{bmatrix} \begin{bmatrix} v_{1} \\ v_{2} \\ v_{3} \\ v_{4} \end{bmatrix} \leq \begin{bmatrix} 0 \\ n \end{bmatrix}$

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Example of controller synthesis

Bounded Producer / Consumer

1) Test

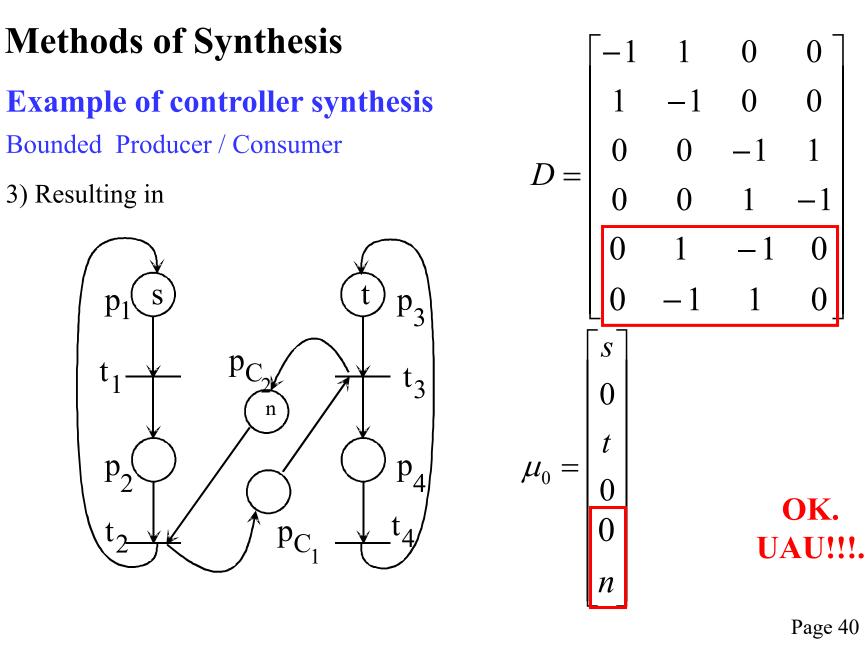
$$b - L\mu_{P_0} = b = \begin{bmatrix} 0\\ n \end{bmatrix} \ge 0.$$
OK.
2) Compute

2) ۰ŀ

$$D_{C}^{-} = \max\left(0, \begin{bmatrix} 0 & -1 & 1 & 0 \\ 0 & 1 & -1 & 0 \end{bmatrix}, 0\right) = \begin{bmatrix} 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix},$$
$$D_{C}^{+} = \max\left(0, 0 - \max\left(0, \begin{bmatrix} 0 & -1 & 1 & 0 \\ 0 & 1 & -1 & 0 \end{bmatrix}\right)\right) - \min\left(0, \begin{bmatrix} 0 & -1 & 1 & 0 \\ 0 & 1 & -1 & 0 \end{bmatrix}\right)$$
$$= \begin{bmatrix} 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix} - \begin{bmatrix} 0 & -1 & 0 & 0 \\ 0 & 0 & -1 & 0 \end{bmatrix} = \begin{bmatrix} 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \end{bmatrix}.$$

and

$$\mu_{C_0} = b - L\mu_{P_0} = \begin{bmatrix} 0\\ n \end{bmatrix}.$$
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0

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()

adding Uncontrollable and Unobservable transitions

Definition of Uncontrollable Transition:

A transition is uncontrollable if its firing **cannot be inhibited** by an external action (e.g. a supervisory controller).

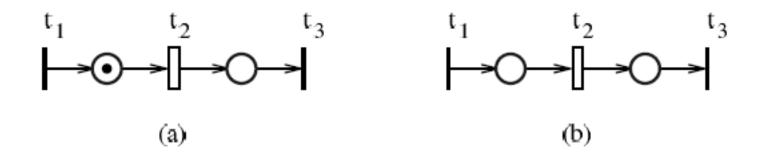
Definition of Unobservable Transition:

A transition is unobservable if its firing **cannot be detected or measured** (therefore the study of any supervisory controller can not depend from that firing).

Proposition:

In a Petri net based controller, both input and output arcs to/from plant transitions are used to trigger state changes in the controller. *Since a controller cannot have arcs connecting to unobservable transitions, then all unobservable transitions are also implicitly uncontrollable*.

adding Uncontrollable and Unobservable transitions



If **t1** is controllable and <u>**t2** is uncontrollable</u>:

- case (a), then t2 cannot be directly inhibited; it will eventually fire
- case (b), then t2 can be indirectly prevented from firing by inhibiting t1.

If <u>t2 is unobservable</u> and t3 is observable, then we cannot detect when t2 fires. The state of a supervisor is not changed by firing t2. However we can **indirectly detect that t2 has fired**, by detecting the firing of t3.

Definition: A marking μ_P is admissible if i) $L\mu_P \leq b$ and ii) $\forall \mu' \in R(C, \mu_P)$ verifies $L\mu' \leq b$

Definition: A Linear Constraint (L, b) is admissible if

i)
$$L\mu_{Po} \leq b$$
 and

ii) $\forall \mu' \in R(C, \mu_{Po})$ such that $L\mu' \leq b$ μ' is an admissible marking.

Note: ii) indicates that the firing of uncontrollable transitions can never lead from a state that satisfies the constraint to a new state that does not satisfy the constraint.

Proposition: Admissibility of a constraint

A linear constraint is admissible *iff*

- The initial markings satisfy the constraint.
- There exists a controller with maximal permissivity that forces the constraint and does not inhibit any uncontrollable transition.

Two sufficient (not necessary) conditions:

Corollary: given a system with uncontrollable transitions, $l^T D_{uc} \le 0$ implies admissibility.

Corollary: given a system with unobservable transitions, $\begin{bmatrix} l^T D_{uo} = 0 \end{bmatrix}$ implies admissibility.

Function MRO_ADM of SPNBOX

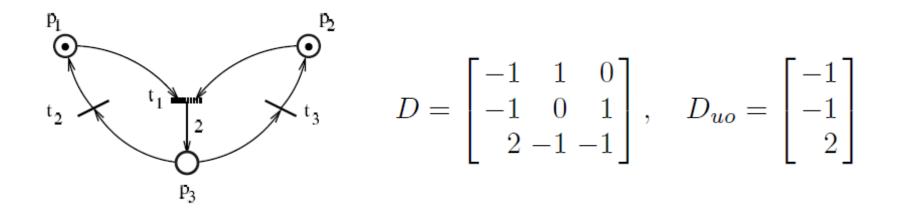
Lemma *: Structure of Constraint transformation

Let $R_1 \in Z^{n_C \times n}$ such that $R_1 \mu_P \ge 0$, $R_2 \in Z^{n_C \times n_C}$ be a matrix with positive elements in the diagonal, If there exists $L' = R_1 + R_2 L$ such that $L' \mu_P \le b'$ $b' = R_2 (b+1) - 1$,

then it is also verified that $L\mu_P \leq b$.

* Lemma 4.10 in [Moody98] pg46

Example: design controller with t1 unobservable (1/4)



Objectives: $\mu_1 + \mu_3 \ge 1$ and $\mu_2 + \mu_3 \ge 1$ which can be written in matrix form as

$$L\mu \le b$$
, $L = \begin{bmatrix} -1 & 0 & -1 \\ 0 & -1 & -1 \end{bmatrix}$, $b = \begin{bmatrix} -1 \\ -1 \end{bmatrix}$

Example extracted from "Supervised Control of Concurrent Systems: A Petri Net Structural Approach", M. Iordache and P. Antsaklis, Birkhauser 2006.

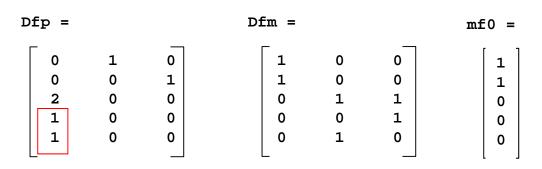
Example: design controller with t1 unobservable (2/4)

- % System and constraints
- $D= \begin{bmatrix} -1 & 1 & 0; \\ -1 & 0 & 1; \\ +2 & -1 & -1]; \end{bmatrix}$
- Dm= -D.*(D<0); Dp= D.*(D>0);

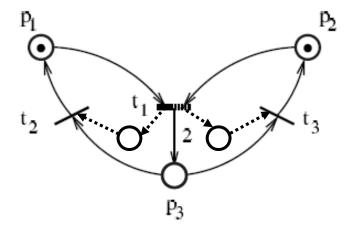
mO= [1 1 0]';

L= -[1 0 1; 0 1 1]; b= -[1; 1];

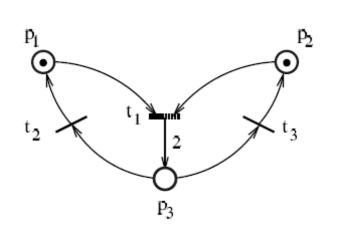
- % Supervisor computation
- [Dfp, Dfm, mf0] =
 linenf(Dp, Dm, L, b, m0);



 $^{\wedge}$ Bad news, supervisor touches t_1 .



Example: design controller with t1 unobservable (3/4)



D= [-1 1 0; -1 0 1; 2 -1 -1]; Tuo= [1]; Tuc= []; L= -[1 0 1; 0 1 1]; b= -[1 1]'; [La, ba, R1, R2] = mro_adm(L, b, D, Tuc, Tuo);

Solution obtained with the function MRO_ADM.m of the SPNBOX toolbox:

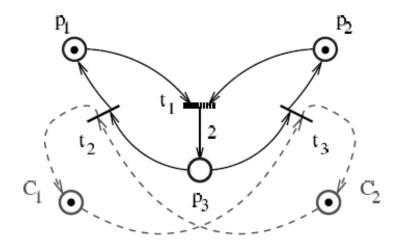
$$R1 = \begin{bmatrix} 0 & 0 & 1 \\ 0 & 0 & 1 \end{bmatrix} \qquad La = \begin{bmatrix} -2 & 0 & -1 \\ 0 & -2 & -1 \end{bmatrix} \qquad ba = \begin{bmatrix} -1 \\ -1 \end{bmatrix}$$
$$R2 = \begin{bmatrix} 2 & 0 \\ 0 & 2 \end{bmatrix} \qquad Note: verify that \quad L_a \mu \le b_a \quad implies \quad L\mu \le b$$

Example: design controller with t1 unobservable (4/4)

Finally the supervised controller is simply obtained from L_a and b_a :

$$\begin{split} D_c &= -L_a D_p \\ &= \begin{bmatrix} -2 & 0 & -1 \\ 0 & -2 & -1 \end{bmatrix} \begin{bmatrix} -1 & 1 & 0 \\ -1 & 0 & 1 \\ 2 & -1 & -1 \end{bmatrix} \\ &= \begin{bmatrix} 0 & 1 & -1 \\ 0 & -1 & 1 \end{bmatrix} \end{split}$$

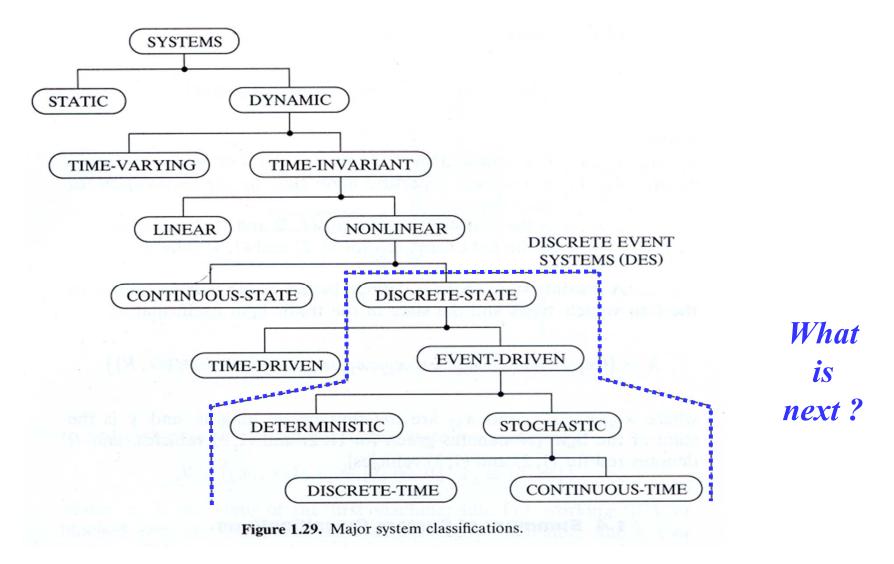
$$\mu_{c0} = b_a - L_a \mu_{p0}$$
$$= \begin{bmatrix} -1 \\ -1 \end{bmatrix} - \begin{bmatrix} -2 & 0 & -1 \\ 0 & -2 & -1 \end{bmatrix} \begin{bmatrix} 1 \\ 1 \\ 0 \end{bmatrix}$$
$$= \begin{bmatrix} 1 \\ 1 \end{bmatrix}$$



Obtained the desired result: supervisor does not touch t_1 .

This course is ending. What is next?

This course is ending.



IST / DEEC / API

Top 10 Challenges in Logic Control for Manufacturing Systems by Dawn Tilbury from University of Michigan

10. Distributed Control	(General management of distributed control applications, Open/distributed control ethernet-based control)
9. Theory	(No well-developed and accepted theory of discrete event control, in contrast to continuous control)
8. Languages	(None of the programming languages do what we need but nobody wants a new programming language)
7. Control logic synthesis	(automatically)
6. Standards	(Machine-control standards every machine is different, Validated standards, Standardizing different types of control logic programming language)
5. Verification	(Standards for validation, Simulation and verification of controllers)
4. Software	(Software re-usability cut and paste, Sophisticated software for logic control, User-unfriendly software)
3. Theory/Practice Gap	(Bridging the gap between industry and academia, Gap between commercial software and academic research)
2. Education	(Educating students for various PLCs, Education and keeping current with evolution of new control technologies, Education of engineers in logic control, Lack of curriculum in discrete-event systems) And the number one challenge in logic control for manufacturing systems is
1. Diagnostics	(Integrating diagnostic tools in logic control, Standardized methodologies for design, development, and implementation of diagnostics)

The End.