

Proof sketches

Asynchronous algorithms: analysis

Lemma A:

Let $\{x(k)\}_{k \in \mathbb{N}}$ sequence of points produced by the asynchronous algorithms I or II

then,

- 1) $\sum_{k \geq 1} \|\nabla \hat{f}(x(k))\|^2 < \infty, a.s.;$
- 2) $\nabla \hat{f}(x(k)) \rightarrow 0, a.s.$

Asynchronous algorithms: analysis

Lemma B:

Let $\{x(k)\}_{k \in \mathbb{N}}$ sequence generated according to Lemma A, with probability one then,

$$\hat{f}(x(k)) \downarrow \hat{f}^*$$

and there exists a subsequence converging to a point in the solution set:

$$x(k_l) \rightarrow y, y \in \mathcal{X}^*$$

Proof of almost sure convergence

Suppose $d_{\mathcal{X}^*}(x(k)) \not\rightarrow 0$

Then, there is an $\epsilon > 0$ and a subsequence $\{x(k_l)\}_{l \in \mathbb{N}}$ such that $d_{\mathcal{X}^*}(x(k_l)) > \epsilon$

As the function is coercive, continuous, and convex, and whose gradient by Lemma A vanishes, then, by Lemma B there is a subsequence of $\{x(k_l)\}_{l \in \mathbb{N}}$ converging to a point in \mathcal{X}^*

Sketch of proof for the almost sure convergence to a point

Fix an $x^* \in \mathcal{X}^*$

Firstly, we prove $\{\|x(k) - x^*\|^2\}_{k \in \mathbb{N}}$ is convergent.

$$\mathbb{E} [\|x(k) - x^*\|^2 | \mathcal{F}_{k-1}] = \sum_{i=1}^n \frac{1}{n} \left\| x(k-1) - \frac{1}{L_{\hat{f}}} g_i(k-1) - x^* \right\|^2$$

$$\|x(k-1) - x^*\|^2 + \frac{1}{nL_{\hat{f}}^2} \left\| \nabla \hat{f}(x(k-1)) \right\|^2 - \frac{2}{nL_{\hat{f}}} (x(k-1) - x^*)^\top \nabla \hat{f}(x(k-1))$$

Sketch of proof for the almost sure convergence to a point

$$(x(k-1) - x^*)^\top \nabla \hat{f}(x(k-1)) = (x(k-1) - x^*)^\top (\nabla \hat{f}(x(k-1)) - \nabla \hat{f}(x^*)) \geq 0$$

$$\mathbb{E} [\|x(k) - x^*\|^2 | \mathcal{F}_{k-1}] \leq \|x(k-1) - x^*\|^2 + \frac{1}{nL_{\hat{f}}^2} \left\| \nabla \hat{f}(x(k-1)) \right\|^2$$

As proved in Lemma A, the sum $\sum_{k \geq 1} \|\nabla \hat{f}(x(k))\|^2 < \infty, a.s.$

so we can invoke the result in Robbins, 1985 to state the convergence of the squared distance.